

May 6, 1973. Stable bundles on a Riemann surface

Sheshadri-Narasimhan thm: Let  $\tilde{C} \rightarrow C$  be the universal covering &  $\Gamma$  the Galois group of a curve  $C$  over  $\mathbb{C}$ . Then if  $V$  is a  $n$ -<sup>unitary</sup> representation of  $\Gamma$  over  $\mathbb{C}$ , the bundle  $\tilde{C} \times^\Gamma V \rightarrow C$

is a direct sum of stable bundles of degree 0. Conversely a vector bundle which is a direct sum of stable bundles of degree zero is obtained in this way.

Weil thm. A bundle  $E$  on  $C$  is obtained from a finite dimensional representation of  $\Gamma \iff$  each of its indecomposable components has degree zero.

Note that  $E$  is obtained from a representation of  $\Gamma \iff$  it has a integrable connection. Since we are over a curve integrability is immediate, so we need only a connection, i.e. a splitting of the sequence

$$(*) \quad 0 \rightarrow E \otimes \Omega \rightarrow J_1(E) \rightarrow E \rightarrow 0.$$

Now  $\text{Ext}^1(E, E \otimes \Omega) = H^1(\text{Hom}(E, E) \otimes \Omega)$

is dual to

$$\begin{aligned} H^0(\text{Hom}(E, E)^\vee) &= \text{Hom}(\text{Hom}(E, E), \mathbb{C}) \\ &= H^0(\text{Hom}(E, E)) \quad \text{by } (\text{trace } AB) \\ &= \text{End}(E, E) \end{aligned}$$

The question now is just what the linear function on

$\text{End}(E)$  belonging to  $(*)$  is. We want to split

$$(1) \quad 0 \longrightarrow \underline{\text{Hom}}(E, E) \otimes \Omega \longrightarrow \underline{\text{Hom}}_{\text{Id}_E}(E, \mathcal{I}_1(E)) \longrightarrow \mathcal{O} \longrightarrow 0$$

which is classified by a class in  $H^1(\underline{\text{Hom}}(E, E) \otimes \Omega)$  dual to  $\text{End}(E)$ .

If  $E$  is ~~a~~ simple, i.e.  $\text{End}(E) = k$ , then

$$H^1(\underline{\text{Hom}}(E, E) \otimes \Omega) \xrightarrow{\text{trace}} H^1(\Omega) \rightarrow 0$$

~~is isomorphic to~~  $\parallel$   
 $k$

both are 1-dimensional and the image of the canonical class is the Atiyah first Chern class, hence it is the degree of  $E$ . Conclude:

Prop: If  $\text{End}(E) = k$ , then  $E$  has a connection (necessarily integrable) iff  $\text{deg}(E) = 0$ .

But in general we have the exact sequence

$$0 \longrightarrow H^0(\mathfrak{sl}(E) \otimes \Omega) \longrightarrow H^0(\mathfrak{gl}(E) \otimes \Omega) \longrightarrow H^0(\Omega) \longrightarrow 0$$

$$\hookrightarrow H^1(\mathfrak{sl}(E) \otimes \Omega) \longrightarrow H^1(\mathfrak{gl}(E) \otimes \Omega) \longrightarrow H^1(\Omega) \longrightarrow 0$$

$\parallel$   
 $k$

where  $\mathfrak{gl}(E) = \underline{\text{End}}(E)$ , and  $\mathfrak{sl}(E)$  is the subsheaf of endos. of trace zero. This exact sequence splits (char. 0). If Weil's theorem is true one has to be able to see why

the class of (1) dies if its ?

The point is that ~~for~~ <sup>given</sup>  $E$  there is a canonical map  $\text{End}(E) \rightarrow k$ , defined by

$$H^0(\text{End}(E)) \times H^1(\text{End}(E) \otimes \Omega) \rightarrow H^1(\Omega)$$

$\alpha$  below  $\mapsto \int \text{tr}(\alpha \cdot \beta) \omega$

and I want a ~~simple~~ formula for this map. The conjecture would be that if

$$E = \bigoplus E_i$$

is a decomposition into indecomposables, then the function is

$$(2) \quad \alpha \mapsto \sum \text{tr}(\alpha|_{E_i}) \cdot \text{deg}(E_i)$$

Possibly a way to proceed is to use the fact that  $E$  has a reduction to the Borel subgroup:

Given  $\alpha \in \text{End}(E)$  we can find a <sup>full</sup> flag in  $E$  which is invariant under  $\alpha$  (do so at the generic point). Then the bundle is reduced to the Borel subgroup and so we have a map of Atiyah extensions

$$\begin{array}{ccccccc}
 0 & \rightarrow & \text{bor}(E) & \longrightarrow & & \longrightarrow & 0 & \longrightarrow & 0 \\
 & & \downarrow & & & & \parallel & & \\
 0 & \rightarrow & \mathfrak{gl}(E) & \longrightarrow & & \longrightarrow & 0 & \longrightarrow & 0
 \end{array}$$

On the other hand  $\alpha \in H^0(\text{bor}(E)) \subset H^0(\mathfrak{gl}(E))$ , so

under the duality pairing  $(A, B) \mapsto \text{tr}(AB)$ ,  $\text{Hom}(E \otimes \mathcal{O}_E \rightarrow \mathcal{O})$   
we just restrict to the Borel subgroup. What this  
clearly gives is the formula

$$\alpha \mapsto \sum \text{tr}(\alpha|L_i) \cdot \text{deg}(L_i)$$

(The last factor results from ~~the~~ the line bundle situation).  
~~Now if E is indecomposable~~

Now if  $E$  is indecomposable<sup>\*</sup> the eigenvalues  
 $\text{tr}(\alpha|L_i)$  must all be the same, so we get

$$\alpha \mapsto \text{tr}(\alpha|E) \text{deg} E$$

\* Construct idempotents from min. poly of  $\alpha$

for an indecomposable bundle, proving (2). Thus  
~~we have~~ we have a proof of Weil's theorem.

Weil's thm: A vector bundle  $E$  on a curve  $C$  has  
a connection (necessarily integrable)  $\iff$  each indecomposable  
factor has degree zero.

# Mumford's method for constructing the moduli space.

Let  $\mathcal{F}$  be a limited family of vector bundles with the same rank and degree, hence the same Hilbert polynomial. Twisting with  $\mathcal{O}(m)$  for  $m$  sufficiently large, one can make  $\dim H^0(E(m))$  independent of  $E$  and such that it generates  $E(m)$ . (More precisely, one arranges that  $H^1(E(m-1)) = 0$  and this implies  $H^1(E(m)) = 0$  and  $H^0(E(m))$  generates  $E(m)$ .) Thus upon choosing a basis for  $H^0(E(m))$  we get a point ~~in~~ in a Hilbert scheme of a certain type. ~~Thus~~ Thus by rigidifying  $E$  to  $(E + V \rightarrow H^0(E(m)))$  we <sup>can</sup> parameterize by a variety of some sort. Now we have to form a quotient by the action of  $\text{Aut}(V)$  to parameterize ~~the~~ the iso. classes of  $\mathcal{F}$ .

~~Now Mumford's idea is to forget the Hilbert scheme~~

In order for Mumford to apply his theory he needs to determine the stable points of the Hilbert scheme. For this he chooses a large number of points on the surface and associates to a vector bundle quotient of  $\mathcal{O} \otimes V$  the fibres at each of the points, thus getting a sequence of quotients  $V \rightarrow W_1, V \rightarrow W_2, \dots, V \rightarrow W_N$  of  $V$ . The idea now is simply to show that if the points are properly chosen then a quotient of  $\text{Aut}(V)$  is stable iff the sequence of points in the Grassmannian of  $V$  is.

Mumford-Sheshadri example of an exact category.

$N$  fixed integer  $\geq 0$ . Objects consist of  $(V, W_1, \dots, W_N)$  where  $V$  is a vector space and  $W_i$  are a sequence of subspaces of the same dimension. An exact sequence

$$0 \rightarrow (V', W_i') \rightarrow (V, W_i) \rightarrow (V'', W_i'') \rightarrow 0$$

is defined to be an exact sequence

$$0 \rightarrow V' \rightarrow V \rightarrow V'' \rightarrow 0$$

such that for each  $i$  we have a map of exact seq.

$$0 \rightarrow W_i' \rightarrow W_i \rightarrow W_i'' \rightarrow 0$$

$$0 \rightarrow \bigwedge V' \rightarrow \bigwedge V \rightarrow \bigwedge V'' \rightarrow 0$$

Call this exact category  $\mathcal{G}^N$ . Here's how it arises:

Let  $X$  be a ~~smooth~~ projective variety over  $k$  (alg. closed say) and  $\mathcal{O}(1)$  a very ample line bundle. Consider the exact category of ~~smooth~~ vector bundles over  $X$  such that  $E$  is generated by  $H^0(E)$  and  $H^1(E) = 0$ . Let  $x_1, \dots, x_N$  be a sequence of closed points. Then

$$E \mapsto (H^0(E), \text{Ker}\{e_{x_i}: H^0(E) \rightarrow E(x_i)\})$$

is exact ~~category~~ from such vector bundles to  $\mathcal{G}^N$ .

$\mathcal{G}^{N'}$  I want to determine the K-theory of  $\mathcal{G}^N$ . Let  $\mathcal{G}^{N'} \subset \mathcal{G}^N$  be the full subcategory consisting of  $(V, W_i)$  such that the  $W_i$  are independent, i.e.  $\bigoplus W_i \subset V$ .

This is closed under extensions since if we have an exact sequence

$$0 \rightarrow (V', W'_i) \rightarrow (V, W_i) \rightarrow (V'', W''_i) \rightarrow 0$$

then we have a map of exact sequences

$$\begin{array}{ccccccc} 0 & \rightarrow & \bigoplus W'_i & \rightarrow & \bigoplus W_i & \rightarrow & \bigoplus W''_i \rightarrow 0 \\ & & \downarrow & & \downarrow & & \downarrow \\ 0 & \rightarrow & V' & \rightarrow & V & \rightarrow & V'' \rightarrow 0 \end{array}$$

so by the serpent lemma  $(V, W_i) \in \mathcal{G}^{N'} \Rightarrow (V', W'_i) \in \mathcal{G}^{N'}$  and  $\mathcal{G}^{N'}$  is closed under extensions. On the other hand we know that given  $(V, W_i) \in \mathcal{G}^N$ , then set

$$\begin{aligned} \bar{V} &= \bigoplus_{j=1}^N V \\ \bar{W}_i &= \bigoplus_{j=1}^N \begin{cases} 0 & j \neq i \\ W_i & j = i \end{cases} \end{aligned}$$

so that we have a map

$$\bar{V} \twoheadrightarrow V$$

carrying  
exact sequence

$$\bar{W}_i \xrightarrow{\sim} W_i. \quad \text{Thus we have an}$$

$$0 \rightarrow (\text{Ker}(\bar{V} \rightarrow V), 0_i) \rightarrow (\bar{V}, \bar{W}_i) \rightarrow (V, W_i) \rightarrow 0$$

showing that any member of  $\mathcal{G}^N$  is a quotient of some member of  $\mathcal{G}^{N'}$ . So the resolution thm. implies  $\mathcal{G}^{N'}$

and  $\mathcal{G}^N$  have the same K-theory.

But now it is clear that  $\mathcal{G}^N$  is the exact category of exact sequences

$$0 \rightarrow \bigoplus W_i \rightarrow V \rightarrow V'' \rightarrow 0$$

so there are obvious characteristic filtrations which show that the K-theory obtained is the direct sum of copies of  $K_*(k)$ ,  $N+1$  copies.

Nothing new is obtained in this way about the K-theory of the curve  $C$ . Thus we get from sending

$$E \mapsto (H^0(E), \text{Ker}\{e\sigma_x : H^0(E) \rightarrow E_x\})$$

the two ingredients of  $K_*(C)$  which we knew about before:

$$\begin{array}{ccc}
K_*(C) & \xrightarrow[\text{to } k]{\text{nest}} & K_*(k) \\
& \searrow & \nearrow \\
& \int_C = H^0 - H^1 & 
\end{array}$$

Question: Describe the exact functors from suff. pos. vector bundles on  $X$  to  $k$ -modules?



## Vector bundles on elliptic curves. (Atiyah's paper)

First let  $E$  be a rank  $n$  vector bundle over  $C$  <sup>a curve</sup> and let  $0 \subset E_1 \subset \dots \subset E_n = E$  be a maximal flag with quotients  $L_i$ . Then the first estimate ~~is~~ is an upper bound for  $\deg(L_i) - \deg(L_{i-1})$ . To derive it we can suppose  $E$  two dimensional

$$0 \rightarrow L_1 \rightarrow E \rightarrow L_2 \rightarrow 0$$

and that

$$0 < \deg(E) + 2(1-g) \leq 2$$

since we can  $\otimes$  with a line bundle. Then RR  $\Rightarrow H^0(E) \neq 0 \Rightarrow \deg(L_1) \geq 0$  so

$$\deg(L_2) - \deg(L_1) \leq \deg(L_2) + \deg(L_1) = \deg E \leq 2g$$

~~Therefore~~

Thus for a maximal flag

$$\deg(L_i) - \deg(L_{i-1}) \leq 2g$$

The next estimate concerns the case where  $E$  is indecomposable whence for  $0 < i < n$

$$0 \neq \text{Ext}^1(E_n/E_i, E_i) \quad \text{is} \quad H^1(\underline{\text{Hom}}(E_n/E_i, E_i))$$

dual to  $H^0(\underline{\text{Hom}}(E_i, E_n/E_i) \otimes \Omega) \neq 0$ . Then fixing  $i$  we have

$$f_i: E_i \xrightarrow{\neq 0} E_n/E_i \otimes \Omega$$

Then let  $j$  be  $\text{minimal}$   $f_i|_{E_j}$  is  $\neq 0$ , whence  $f_i: L_j \hookrightarrow E_n/E_i \otimes \Omega$

and so  $\deg(L_j) - (2g-2) \leq \deg(L_{i+1})$

for some  $1 \leq j \leq i$ . Thus want

$$\deg(L_i) \geq \deg(L_1) - (i-1)(2g-2)$$

E ind.

and we can get this by induction

$$\begin{aligned} \deg(L_{i+1}) &\geq \deg(L_j) - (2g-2) \\ &\geq \deg(L_1) - (j-1)(2g-2) - (2g-2) \\ &\geq \deg(L_1) - (i)(2g-2) \quad \text{since } j \leq i. \end{aligned}$$

But now suppose we have an elliptic curve whence  $\Omega = \mathcal{O}$ . Here to show  $E_i$  can be chosen so that  $L_1 \subset L_i$  for all  $i$ . Now again using induction we have

$$f_i : E_i \xrightarrow{\ast} E_n / E_i$$

hence  $\exists j \leq i$

$$f_i : L_j \hookrightarrow E_n / E_i$$

~~so we can suppose  $L_{i+1}$  chosen  $\dots$~~  If the maximal degree =  $\deg(L_{i+1})$  is the same as  $\deg(L_j)$  we can suppose  $L_{i+1}$  chosen to be  $f_i(L_j)$ . Otherwise

$$\deg(L_j) < \deg(L_{i+1})$$

and we know  $\exists L_j \hookrightarrow L_{i+1}$ . In either case  $L_{i+1}$  can be assumed to contain  $L_{01}$

Now suppose  $E$  is indecomposable of rank  $r$  and

$\deg(E) = d$  where  $0 < d \leq n$ . ~~Assume that  $d > n$~~

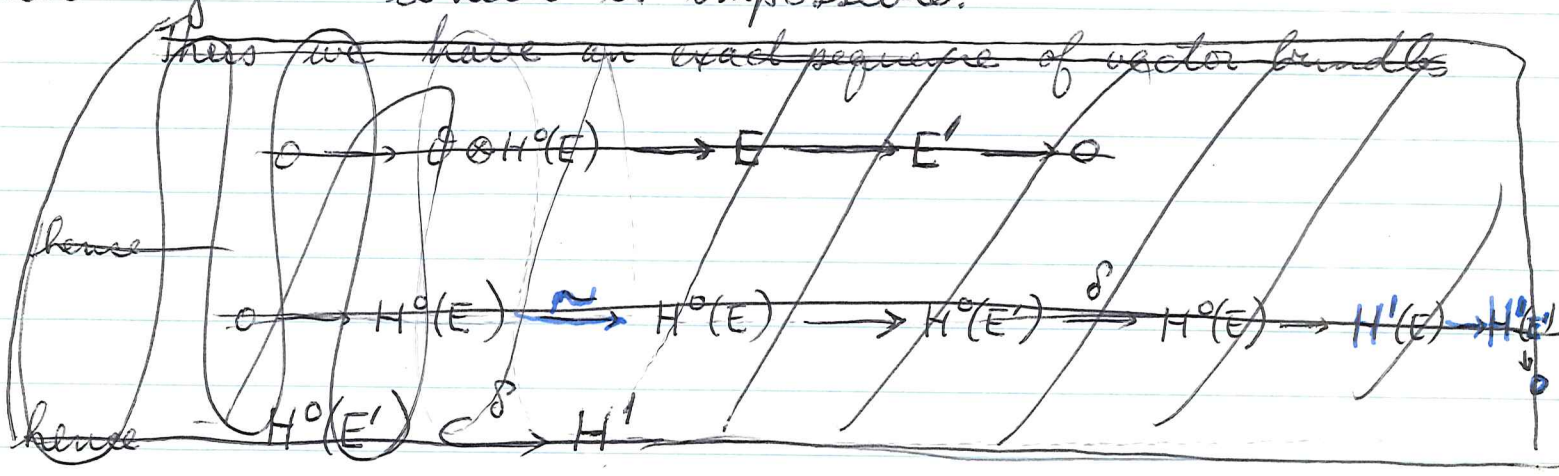
~~then there is no subbundle of  $E$  of degree  $> d$ .~~  
 We know that  $H^0(E) \neq 0$  by RR, hence  $\deg(L_1) \geq 0$ .

If  $\deg(L_1) = 1$ , then ~~then~~  $L_1 \subset L_i$  &  $\sum \deg L_i = d \leq n$   
 $\Rightarrow d = n$  and all  $L_i = L$ , so  $E = L_1 \otimes F_n$  where  $F_n$  is the unique indecomposable bundle which is a successive extension of  $\mathcal{O}$ 's of rank  $n$ . Thus  $E$  is semi-stable of slope 1.

On the other hand suppose  $\deg(L_1) = 0$ . ~~I~~ I claim that the map

$$\mathcal{O} \otimes H^0(E) \longrightarrow E$$

is a vector bundle injection. To see this it is enough to show that  $s(x) \neq 0 \quad \forall x \in X, \forall s \in H^0(E), s \neq 0$ . But otherwise the subline bundle  $L$  of  $E$  generated by  $s$  would have  $\deg > 0$  which is impossible.



But  $h^0(E) \geq \deg(E)$ , so that the first thing to note is that  $d \geq n$  is impossible in this case since otherwise  $E$  would be trivial.

If  $d = n - 1$ , then must have  $h^0(E) = n - 1$ , so

$E$  is an extension of  $\Lambda^n E$  by  $\mathcal{O}^{n-1}$ . In general  
~~define~~ define  $E'$  by

$$0 \rightarrow \mathcal{O} \otimes H^0(E) \rightarrow \del{E} E \rightarrow E' \rightarrow 0$$

so that  $E' = \Lambda^n E$  if  $d = n-1$ . Since  $h^0(E) - h^1(E) = n-1$   
 and  $h^0(E) = n-1$  we have  $h^1(E) = 0$  and therefore  
 the above should be the ~~canonical~~ canonical extension?

Better: Let  $L$  have degree  $k$ . Then

$$\begin{aligned} \text{Ext}^1(L, \mathcal{O} \otimes_r V) & \del{\text{is isomorphic to } H^1(L) \otimes_r V} \\ & = H^1(L) \otimes_r V = \text{Hom}(H^0(L), V) \end{aligned}$$

so there is a canonical extension

$$0 \rightarrow \mathcal{O} \otimes H^0(L) \rightarrow ? \rightarrow L \rightarrow 0$$

and from the coh. exact sequence we get

$$0 \rightarrow H^0(L) \rightarrow H^1(\mathcal{O} \otimes H^0(L)) \rightarrow H^1(?) \rightarrow 0$$

hence  $H^1(?) = 0$ . It is clear that  $?$  is indecomposable  
 since any endo. of  $?$  induces one of the exact sequences.

Suppose  $d = n-2$ .  $h^0(E) - h^1(E) = d$ . There  
 are two possibilities:  $h^0(E) = n-1$  so  $E'$  is a line bundle  
 of degree  $n-2$

$$\begin{array}{ccccccc} 0 & \rightarrow & \mathcal{O}^{n-1} & \rightarrow & E & \rightarrow & L \rightarrow 0 \\ & & \uparrow & & \uparrow & & \parallel \\ 0 & \rightarrow & \mathcal{O} \otimes H^0(L) & \rightarrow & ? & \rightarrow & L \rightarrow 0 \end{array}$$

But if  $H^0(L)$  is of dim  $n-2$ , then the image of  $\mathcal{O} \otimes H^0(L)$  sits is a proper ~~sub~~ bundle of  $\mathcal{O}^{n-2}$  and so it is clear that  $E = ? \oplus \mathcal{O}$ . This case doesn't happen.

Thus  $h^0(E) = n-2$ ,  $h^1(E) = 0$  and  $E'$  is of rank ~~two~~ two and  $E$  is the canonical extension

$$0 \rightarrow \mathcal{O} \otimes H^0(E) \rightarrow ? \rightarrow E' \rightarrow 0$$

It follows that  $E'$  is indecomposable of rank 2 and degree  $n-2$ .

In general let  $E$  be indecomposable with degree  $d > 0$  and  $\deg(L_1) = 0$  so that we have

$$\begin{array}{ccccccc} 0 & \rightarrow & \mathcal{O} \otimes H^0(E) & \rightarrow & E & \rightarrow & E' \rightarrow 0 \\ & & \uparrow \alpha & & \uparrow & & \parallel \\ 0 & \rightarrow & \mathcal{O} \otimes H^0(E') & \rightarrow & ? & \rightarrow & E' \rightarrow 0 \end{array}$$

where the bottom is the canonical extension. Since  $H^0(E') \xrightarrow{\cong} H^0(\mathcal{O} \otimes H^0(E)) \cong H^0(E)$ , the map  $\alpha$  if it is onto then it is an isom. But if  $\alpha$  isn't onto then we get

$$0 \rightarrow ? / \mathcal{O} \otimes \ker \alpha \rightarrow E \xleftarrow{\cong} \mathcal{O} \otimes \text{Coker } \alpha \rightarrow 0$$

so  $\alpha$  has to be an isomorphism. Conclude that ~~any~~ an indecomposable with  $\deg = d > 0$  and  $\deg(L_1) = 0$  is the canonical extension of ~~an~~ an  $E'$ . Clearly  $E'$  is indecomposable of the same degree as  $E$ . ~~But~~

~~$H^0(E)$~~

Lemma: Let  $E$  a vector bundle with  $\deg(L_1) = 0$ .  
Then the map  $\mathcal{O} \otimes H^0(E) \rightarrow E$   
is a subbundle inclusion.

L This holds for any curve.

For an elliptic curve with  $E$  decomposable we have putting  $E' = E / \mathcal{O} \otimes H^0(E)$

~~$H^0(E)$~~

$$\begin{aligned} \text{Ext}^1(E', \mathcal{O} \otimes_k V) &= H^1(E'^*) \otimes V \\ &= H^0(E' \otimes \Omega)^* \otimes V \\ &= \text{Hom}(H^0(E' \otimes \Omega), V) \end{aligned}$$

so we have a canonical extension

$$\begin{array}{ccccccc} 0 \rightarrow \mathcal{O} \otimes H^0(E' \otimes \Omega) & \rightarrow & \tilde{E}'_{\text{can}} & \rightarrow & E' & \rightarrow & 0 \\ & \uparrow \alpha & \downarrow & & \parallel & & \leftarrow \text{hence a map} \\ 0 \rightarrow \mathcal{O} \otimes H^0(E) & \rightarrow & E & \rightarrow & E' & \rightarrow & 0 \end{array}$$

and when  $E$  is indecomposable it follows that  $\alpha$  is onto. But  $\delta: H^0(E') \hookrightarrow H^1(\mathcal{O} \otimes H^0(E)) \cong H^0(E)$ , hence  $\alpha$  has to be ~~onto~~ an isomorphism. Clearly  $E'$  is indecomposable.

Conversely suppose  $E'$  indecomposable and consider the canonical extension:

$$0 \rightarrow \mathcal{O} \otimes H^0(E' \otimes \Omega) \rightarrow \tilde{E}'_{\text{can}} \rightarrow E' \rightarrow 0$$

It's pretty clear that  $\rho: H^0(E') \xrightarrow{\sim} H^1(\mathcal{O}) \otimes H^0(E' \otimes \Omega)$  so that  $H^0(E' \otimes \Omega) = H^0(\tilde{E}'_{\text{can}})$ . In particular any endo. of  $E'_{\text{can}}$  induces one of the whole exact sequence, and so  $\tilde{E}'_{\text{can}}$  has to be indecomposable when  $E'$  is

So now start with an indecomposable  $E_n$  <sup>of deg > 0</sup> such that  $\deg(L_1) = 0$  and form  $E'$  and note that  $H^0(E) = H^0(E')$  and  $H^1(E) = H^1(E')$ . Now look at  $\deg(L'_i)$  for  $E'$ . If  $> 0$ , then  $H^1(E') = 0$  since  $L'_i \subset L_i$  for all  $i$ . So by induction we can establish that  $H^1(E) = 0$ , ~~so~~

~~rank(E') = rank(E) - d~~  $h^0(E) = d.$

Now ~~rank(E')~~ note that if we start with ~~deg~~  $0 < d < n$  then  $\deg(L_1) = 0$ , and so

$$\text{rank}(E') = \text{rank}(E) - d$$

$$\deg(E') = d.$$

On the other hand, ~~deg~~ once we arrange  $0 < \deg E \leq \text{rg}(E)$  then  $\deg(L_1) = 0$  ~~iff~~  $\deg(E) < \text{rg}(E)$ . (Recall:  $\deg(L_1) > 0 \Rightarrow$  (as  $L_1 \subset L_i$ ) all  $L_i = L_1$  so  $\deg(E) = \text{rg}(E)$ . But also if  $\deg(L_1) = 0$ , then  $\mathcal{O} \otimes H^0(E) \hookrightarrow E + h^0(E) \geq \text{rg}(E)$  impossible)  
Thus for  $n > d$  we have

$$E(n, d) \simeq E(n-d, d)$$

$$E \longmapsto E'$$

where  $E(n, d)$  denotes the set of iso. classes of indecomposable

bundles of rank  $n$  and degree  $d$ .

Thm: All  $\mathcal{E}(n, d)$  are in 1-1 correspondence with the elliptic curve  $C$ .

Proof: <sup>Use induction on  $n$</sup>  By tensoring with a line bundle of deg  $k$  get

$$\mathcal{E}(n, d) \simeq \mathcal{E}(n, d+kn)$$

so can suppose  $0 < d \leq n$ . If  $d=n$  we know that we have an isom with  $C: L \mapsto L \otimes F_n$ . If  $d < n$ , then we have (as above) an isomorphism

$$\mathcal{E}(n, d) \xrightarrow{\sim} \mathcal{E}(n-d, d)$$

so the induction marches.

Prop:  $E$  indecomposable of degree  $> 0 \Rightarrow H^1(E) = 0$ .

Pf: If  $\deg(L_i) \geq 0$ , then because  $L_1 \subset L_2 \subset \dots \subset E$  is a successive extension of line bundles of  $\deg > 0$ , hence  $H^1(E) = 0$ . On the other hand if  $\deg(L_i) = 0$ , then we have  $H^1(E) = H^1(E')$  where  $E'$  is indecomposable, + of the same degree, but with smaller rank.

When rank  $E \geq 2$ ,

Prop: An indecomposable  $E$  such that  $\frac{\deg E}{\text{rank } E} \notin \mathbb{Z}$  is stable and conversely.

Proof: Any stable bundle is indecomposable (in fact  $\text{End}(E) = k$ ), and we have seen that if  $\text{rank } E$  divides  $\deg E$ ,



Proposition: Let  $0 \rightarrow E' \rightarrow E \rightarrow E'' \rightarrow 0$  be an exact sequence of vector bundles over a curve  $C$  of genus  $g$ . If the sequence doesn't split, then

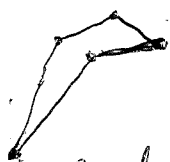
$$\mu_{\min}(E') - \mu_{\max}(E'') \leq 2g - 2.$$

Proof:  $\text{Ext}^1(E'', E') = H^1(\underline{\text{Hom}}(E'', E'))$  is dual to  $H^0(\underline{\text{Hom}}(E', E'') \otimes \Omega) = \text{Hom}(E', E'' \otimes \Omega)$ , so if the sequence doesn't split  $\exists$

$$f: E' \rightarrow E'' \otimes \Omega \neq 0.$$

Then  $\text{Coim}(f) = E' / \text{Ker} f$  is a  $\neq 0$  quotient bundle of  $E'$  so

$$\mu_{\min}(E') \leq \mu(\text{Coim} f)$$



and  $\text{Im}(f)$  is a  $\neq 0$  subbundle of  $E'' \otimes \Omega$  so

$$\mu(\text{Im}(f)) \leq \mu_{\min}(E'' \otimes \Omega) = \mu_{\max}(E'') + 2g - 2$$

and since  $\text{Coim} f \hookrightarrow \text{Im} f$  have the same rank

$$\mu(\text{Coim} f) \leq \mu(\text{Im} f).$$

The proposition follows.

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Corollary: On an elliptic curve an indecomposable bundle is semi-stable.

In effect there can be only one slope

it is indecomposable and  
 (Prop. On an elliptic curve a bundle is stable iff its rank and degree are relatively prime ( $\mathbb{Z}d + \mathbb{Z}r = \mathbb{Z}$ ).

Quasi-Proof:  $E$  indecomposable  $\Rightarrow E$  semi-stable. If  $(d, r) = 1$ , then have for any  $E' \subset E$  that  $\frac{d'}{r'} \leq \frac{d}{r}$ ; must be  $<$  ~~because~~ because by assumption  $\frac{d}{r}$  is in lowest terms.

For the converse suppose  $E$  is stable of degree  $d$  rank  $r$  with  $(d, r) = 1$ , and consider the bundle  $E \otimes F_n$  with invariants  $(dn, rn)$ . Then  $E \otimes F_n$  is semi-stable with quotients isom to  $E$ .  
 ~~$E \otimes F_n$~~

$$0 \subset E \otimes F_1 \subset E \otimes F_2 \subset \dots \subset E \otimes F_n$$

Claim this filtration is characteristic; by induction one should be able to see that

$$\text{Hom}(E, E \otimes F_n) = k$$

$$\text{Ext}^1(E, E \otimes F_n) = k$$

so that any idempotent operator  <sup>$\neq 0$</sup>  has to preserve the filtration and so  $\bullet = 1$ . Thus  $E \otimes F_n$  has to be indecomposable.

So now more or less it is clear that  $E \mapsto E \otimes F_n$  will be a map from  $\mathcal{E}(d, r) \rightarrow \mathcal{E}(nd, nr)$  which has to be onto by isogeny rules. This point should be in Atiyah's paper.

---

no good

May 10, 1973

Possible proof of  $K_g(A[t]) = K_g(A) \oplus Nil_{g-1}(A)$ :

Try to lift  $A[t]$ -modules to graded  $A[T_0, T_1]$ -modules.  
~~Recall that~~ Let  $Pgr(A[T_0, T_1])$  be the cat. of f.g. proj graded  $A[T_0, T_1]$ -modules  $P = \bigoplus_{n \geq 0} P_n$ . Recall that non-canonically

$$(1) \quad P \cong \bigoplus_{k=0}^{\infty} A[T_0, T_1](-k) \otimes_A L_k \quad L_k \in P(A)$$

and that

$$(2) \quad K_g(Pgr(A[T_0, T_1])) \cong K_g(A)[z]$$
$$(A[T_0, T_1](-k) \otimes_A ?)_* \times \longleftarrow \times z^k$$

Now we have the localization map

$$Pgr(A[T_0, T_1]) \longrightarrow P(A[t])$$

$$P \longmapsto P/(T_0 - 1)P$$

$t = \text{image of } T_1$

and it is clear from (1) that the image of this consists of objects of the form  $A[t] \otimes_A L$ ,  $L \in P(A)$ .  
Call this full subcategory  $P' \subset P(A[t])$ .

Now given  $V \in P'$  we have the building consisting of those  $P$  such that

$$P/(T_0 - 1)P \xrightarrow{\sim} V.$$

Better: Form the category  $\mathcal{C}$  whose objects are  $P \in \text{Pgr}(A[T_0, T_1])$  and in which a morphism  $P \rightarrow P'$  is an inclusion whose kernel  $P'/P$ ,  $T_0$  is nilpotent and such that  $P'/P \in \text{Pgr}(A[T_1])$ . ~~Suppose that~~  
 Then have

$$(3) \quad \lambda : \mathcal{C} \longrightarrow \text{Iso}(P')$$

$$P \longmapsto P/(T_0-1)P = \varinjlim \{ P_0 \xrightarrow{T_0} P_1 \xrightarrow{T_0} P_2 \rightarrow \dots \}$$

which we want to show is a homotopy equivalence. So we consider  $\lambda/V$ . ~~Observe that~~ Observe that

$$P_0 \hookrightarrow P_1 \hookrightarrow P_2 \hookrightarrow \dots \varinjlim P_n = P/(T_0-1)P$$

so that ~~given~~ given  $P/(T_0-1)P \cong V$  in  $\lambda/V$  we can identify  $P_n$  with an ~~A~~  $A$ -submodule of  $V$ . In fact to give an  $A[T_0, T_1]$ -module  $M$  which is  $T_0$ -torsion-free and such that  $\lambda M \cong V$  is the same thing as giving ~~a sequence of~~ a sequence of  $A$ -submodules  $M_n \subset V$ ,  $n \geq 0$ , such that  $M_n \subset M_{n+1}$  and  $\lambda M_n \subset M_{n+1}$ , and  $\cup M_n = V$ . For  $M$  to be in  $\text{Pgr}(A[T_0, T_1])$  <sup>probably</sup> means that

$$P_n/P_{n-1} + tP_{n-1} \in \mathcal{P}(A) \quad (=0 \quad n \gg 0)$$

for each  $n$ , and that  $T_1: P_n/P_{n-1} \rightarrow P_{n+1}/P_n$  is injective  $\forall n$ . Thus these conditions imply

$$P/(T_0, T_1)P \in \mathcal{P}(A)$$

$$\text{Tor}_1^{A[T_0, T_1]}(A, P) = 0$$

so by the lemma ~~in the paper~~ in the paper,  $P \in \text{Pgr}(A)$ .

~~So it is clear that  $\lambda/V$  is equivalent to~~

So it is clear that  $\lambda/V$  is equivalent to the ordered set of such families  $\{P_n \subset V\}$ , where  $\{P_n\} \leq \{P'_n\} \iff P_n \subset P'_n$  for all  $n$  and

$$\bigoplus_n P'_n/P_n \in \text{Pgr}(A[T_1])$$

.. is killed by  $T_0^N$

I want to show this ordered set is directed (below). Thus given  $P, P'$  I want to find  $P'' \leq P, P'$ . Since  $P$  is fin. gen.  $T_0^N P \subset P'$  for some  $N$ , and since  $T_0^N P \leq P$  we can suppose  $P \subset P'$ . Now  $\exists n \exists T_0^n P' \subset P$ . Then have

$$0 \rightarrow P''/T_0^n P' \rightarrow P'/T_0^n P' \rightarrow P'/P \rightarrow 0$$

Now because  $P, P'$  are projective over  $A[T_1]$ , we know  $P'/P$  has projective dim 1 over  $A[T_1]$ , hence  $P/T_0^n P'$  is projective over  $A[T_1]$ . ~~But if  $n$  is large enough so that  $T_0^n P \subset P'$ , then~~ But  $P/T_0^n P'$  is a quotient of  $P/T_0^n P$ , hence  $P/T_0^n P$  is finitely gen. over  $A[T_1]$ . Thus  $T_0^n P' \leq P, P'$  as desired.

Let  $\bar{C}$  be the arrow cat. of  $C$ , and let  $\mathcal{H}$  denote the category of those gr  $A[T_0, T_1]$ -modules  $M = M_0 \oplus \dots$  such that (i)  $T_0$  nilp. on  $M$ , (ii)  $M \in \text{P}(A[T_1])$ . Then we have a functor

$$(4) \quad \bar{C} \longrightarrow Q(\mathcal{H}) \quad (P \hookrightarrow P') \mapsto P'/P.$$

Claim this functor is fibred and the fibre over  $M \in \mathcal{K}$  is the groupoid of  $P \twoheadrightarrow M$  with  $P \in \text{Pgr}(A[T_0, T_1])$ .  
 Clearly this is the fibres, so all we have to do is to check ~~that~~ that if we have  $P \twoheadrightarrow M$ , then the kernel is in ~~in~~  $\text{Pgr}(A[T_0, T_1])$ . But ~~we~~ we have the char. sequence

$$0 \longrightarrow A[T_0, T_1] \otimes_{A[T_1]} M \longrightarrow A[T_0, T_1] \otimes_{A[T_1]} M \longrightarrow M \longrightarrow 0$$

showing  $M \in \text{Pgr}(A[T_0, T_1])$ , so this is clear.

May 11

(What I am trying to do is this: Identify  $\text{P}(A[t])$  modules with graded  $A[T_0, T_1, T_0^{-1}]$ -modules and construct the localization theorem in this graded setting)

We now have that  $\mathcal{C} \xrightarrow{\text{les}} \text{Iso}(\mathcal{P}')$  and that  $\bar{\mathcal{C}}$  (= arrow category of  $\mathcal{C}$ ) fibres over  $Q(\mathcal{H})$ . Thus we get a spectral sequence

$$E_{Pq}^2 = H_p(Q(\mathcal{H}), M \mapsto H_q(\{P \twoheadrightarrow M\})) \implies H_q(\text{Iso } \mathcal{P}')$$

~~which~~ which I want to localize.

So the idea is to let  $I = \text{Iso}(\text{Pgr}(A[T_0, T_1]))$  act on the category  $\bar{\mathcal{C}}$  and  $\text{Iso}(\mathcal{P}')$  and we get a fibration

$$\bar{I}_I \longrightarrow \text{Iso } \mathcal{P}'_I \longrightarrow Q(\mathcal{H})$$

and hence a long exact homotopy sequence

$$K_0 \mathcal{H} \xrightarrow{\alpha} K_0(\text{Proj } A[T_0, T_1]) \longrightarrow K_0 A[t] \longrightarrow K_{-1} \mathcal{H} \longrightarrow \dots$$

$$\quad \quad \quad \parallel$$

$$\quad \quad \quad K_0 A[z]$$

Next we have to compute the map  $\alpha$ . ?

---

$A = k$  field,  $B = k[T_0, T_1]$ . Take  $V = k[t]^2$ .  
 Now consider those  $P = P_0 \oplus \dots \oplus P_n \oplus \dots$  free over  $B$   
 of rank 2 such that  $V = \text{im } P / (T_0 - 1)P$ . Such a  
 $P$  is isomorphic to  $B(-p) \oplus B(-q)$  where  $0 \leq p \leq q$ .

June 6, 1973:  $p^n A = 0 \implies K_n A [p^{-1}] \xrightarrow{\sim} K_n(A[t]) [p^{-1}]$ .

We can prove the fundamental thm. as follows

$$\mathbb{P}_A^1 \xleftarrow{f} Sp(A[t^{-1}])$$

$$\partial \rightarrow K_0(\text{Nil } A) \rightarrow K_0(\mathbb{P}_A^1) \xrightarrow{f^*} K_0(A[t^{-1}]) \xrightarrow{\partial} \dots$$

~~Nil~~  $\text{Nil}(A) = A[t]$ -mods.  $M \ni M \in \mathcal{P}(A), t^n M = 0$  some  $n$

Think of these as sheaves on  $\mathbb{P}_A^1$  with support at  $t=0$ . which are finite loc. free over  $A$ .

I want now to consider the effect of a homom  $A[t] \rightarrow A[x]$  sending  $t \mapsto x^n$ .

$$A[T_0, T_1] \rightarrow A[X_0, X_1] \quad T_i \mapsto X_i^n$$
  
$$t = \frac{T_1}{T_0} \mapsto \left(\frac{X_1}{X_0}\right)^n = x^n$$

certainly finite ~~locally~~ locally free map:  $\mathbb{P}_A^1 \xrightarrow{f} \mathbb{P}_A^1$   
 $x^n \leftarrow t$

so it ought to induce a map of localization sequences

$$\begin{array}{ccccc} \text{Nil}(A) & \xrightarrow{\quad} & \mathcal{P}(\mathbb{P}_A^1) & \xrightarrow{f^*} & \mathcal{P}(A[x^{-1}]) \\ \downarrow f_* & & \downarrow f_* & & \downarrow f_* \\ \text{Nil } A & \xrightarrow{\quad} & \mathcal{P}(\mathbb{P}_A^1) & \xrightarrow{f^*} & \mathcal{P}(A[t^{-1}]) \end{array}$$

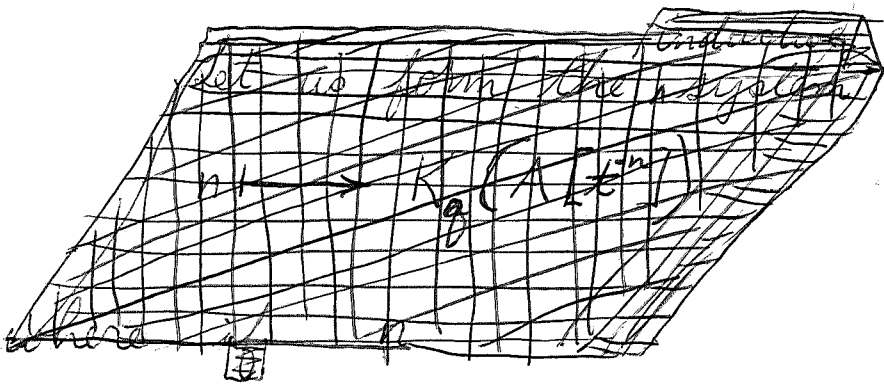
$f_*$  of  $M$  with  $x$  acting as  $\nu: M \rightarrow M$  is  $M$  with  $t$  acting as  $x^n$ .



Can also check ~~this~~ <sup>this</sup> using the map

$$\begin{array}{ccccc} \text{Nil}(A) & \longrightarrow & \mathcal{P}_1(A[x]) & \longrightarrow & \mathcal{P}_1(A[x, x^{-1}]) \\ \downarrow & & \downarrow f_* & & \downarrow f_* \\ \text{Nil}(A) & \longrightarrow & \mathcal{P}_1(A[t]) & \longrightarrow & \mathcal{P}_1(A[t, t^{-1}]) \end{array}$$

So we seem to be able to do the following:



Suppose  $A$  is of characteristic  $p$ . Then I can consider the operation

$$\begin{array}{ccccc} K_0(A[t]) & \xrightarrow{f_*} & K_0(A[t^p]) & \xrightarrow{f^*} & K_0(A[t]) \\ \uparrow i^* & & \uparrow i^* & & \\ K_0 A & \xrightarrow{p} & K_0 A & & \end{array}$$

~~where~~  $f_* i^* = p i^*$  because

$$f_* i^*(E) = A[t] \otimes_A E = \coprod_{i=0}^{p-1} A[t^p] t^i \otimes_A E$$

Now ~~show~~

$$f^* f_* V = (A[t] \otimes_{A[t^p]} A[t]) \otimes_{A[t]} V$$

and because we are in characteristic  $p$  the aug. ideal in  $A[t] \otimes_{A[t^p]} A[t]$  is nilpotent

$$(t \otimes 1 - 1 \otimes t)^p = 0$$

etc., so it should be the case that

$$f^* f_* = \text{mult. by } p.$$

and so iterating

~~$$f^* f_* = p$$~~

$$(f^2)^* (f^2)_* = f^* f^* f_* f_* = p f^* f_* = p^2$$

etc. But now ~~show~~ have

Lemma: ~~Let~~  $\mathcal{P}$  ~~be a set of primes~~  $\text{Nil}^n(A) = A[t]/t^n \text{ mod } \mathcal{P}$  which are in  $\mathcal{P}(A)$ , then

$$K_* \text{Nil}(A) = \varinjlim_n K_* \text{Nil}^n(A)$$

We have computed that

$$\begin{array}{ccc} K_0(A[t]) & \xrightarrow{f_*} & K_0(A[t^p]) \\ \downarrow d & & \downarrow d \\ \text{Nil}_{g-1}(A) & \xrightarrow{f} & \text{Nil}_{g-1}(A) \end{array}$$

commutes where  $g$  sends  $(P, v)$  to  $(P, v^p)$ . Thus any element of  $\text{Nil}_{g-1}(A)$  gets killed by  $(f^N)^*$ , so we get

Prop: If  $A$  of char.  $p$ , then  $\text{Nil}_g(A)$  is  $p$ -torsion for  $g \geq 0$ , hence

$$K_{\frac{g}{p+1}}(A)[p^{-1}] \xrightarrow{\sim} K_{\frac{g}{p+1}}(A[t])[p^{-1}] \text{ for } g \geq 0.$$

But we <sup>also</sup> know that  $K_0 A$  is canonically a direct summand of  $K_1(A[x, x^{-1}])$ . Thus from the isom

$$K_1(A[x, x^{-1}])[p^{-1}] \xrightarrow{\sim} K_1(A[x, x^{-1}][t])[p^{-1}]$$

"  $A[t][x, x^{-1}]$

one sees the above <sup>prop.</sup> holds for  $K_0$  also.

Question: Is ~~any~~  $K_{-n}(A)[\frac{1}{p}] = 0$ ,  $-n < 0$  ?

This amounts to showing that  $K_{-1}(A)[\frac{1}{p}] = 0$

Question: If  $A$  is noetherian of char.  $p$ , is  $K_n A[\frac{1}{p}] \xrightarrow{\sim} K'_n(A)[\frac{1}{p}]$  ?

**NO**, take  $A$  to be affine singular curve, whose ~~Picard group~~ Picard group has a  $\mathbb{G}_m$  part. Precisely take a ~~the~~ plane cubic  $X$  with ordinary



double point; its normalization  $\bar{X}$  is  $\mathbb{P}^1$  (project thru the double point). One has

$$K_0(\bar{X}) = \mathbb{Z} \oplus \text{Pic}(\bar{X}) = \mathbb{Z} \oplus \mathbb{Z}$$

$$K_0(X) = \mathbb{Z} \oplus \text{Pic}(X) \text{ ~~is not equal to } K_0(\bar{X})~~$$

$$0 \longrightarrow k^* \longrightarrow \text{Pic}(X) \longrightarrow \underset{\mathbb{Z}}{\text{Pic}(\bar{X})} \longrightarrow 0$$

e.g. <sup>a</sup> ~~line bundles~~ line bundles on  $X$  is same as one on  $\bar{X}$  with a way of identifying the fibres at  $P$  and  $Q$ . Thus

~~is not equal to~~

$$K_0 X = \mathbb{Z} \oplus (\mathbb{Z} \oplus k^*)$$

~~is not equal to~~

$$\begin{array}{ccccccc}
K_1 F & \longrightarrow & \coprod_{x \in X} \mathbb{Z} & \longrightarrow & K'_0 X & \longrightarrow & K'_0 F \longrightarrow 0 \\
\uparrow s & & \uparrow & & \uparrow & & \uparrow s \\
K_1 F & \longrightarrow & \coprod_{x \in \bar{X}} \mathbb{Z} & \longrightarrow & K'_0 \bar{X} & \longrightarrow & K'_0 F \longrightarrow 0
\end{array}$$

yields

$$\mathbb{Z} \xrightarrow{\text{rg, deg}} K'_0 \bar{X} \xrightarrow{\mathbb{Z} \oplus \mathbb{Z}} K'_0 X \longrightarrow 0$$

$1 \mapsto P-Q$

so

$$K'_0 X = \mathbb{Z} \oplus \mathbb{Z}$$

and hence the map  $K_0 X \rightarrow K'_0 X$  has  $k^*$  for its kernel. This will persist upon removing a simple point from  $X$ , so we can make this example affine if we wish.

Direct proof that  $\text{Nil}_*(A)$  is  $p$ -torsion when  $p \cdot A = 0$ .

Consider  $(M, \nu) \in \text{nil}^{p^2}(A)$  so that  $\nu^{p^2} = 0$ .  
Then we will show that

$$A[t] \otimes_{A[t^{p^2}]} M \quad t^{p^2} \text{ acts as } \nu^{p^2}$$

has a canonical filtration with  $p^2$  quotients all isomorphic to  $M$ .

$$A[t] \otimes_{A[t^{p^2}]} A[t] = A \otimes_{\mathbb{F}_p} \left( \mathbb{F}_p[t] \otimes_{\mathbb{F}_p[t^{p^2}]} \mathbb{F}_p[t] \right)$$

$$I = (t \otimes 1 - 1 \otimes t) \text{ in } \mathbb{F}_p[t] \otimes_{\mathbb{F}_p[t^{p^2}]} \mathbb{F}_p[t] = \mathbb{F}_p[t, t''] / (t'^{p^2} - t''^{p^2})$$

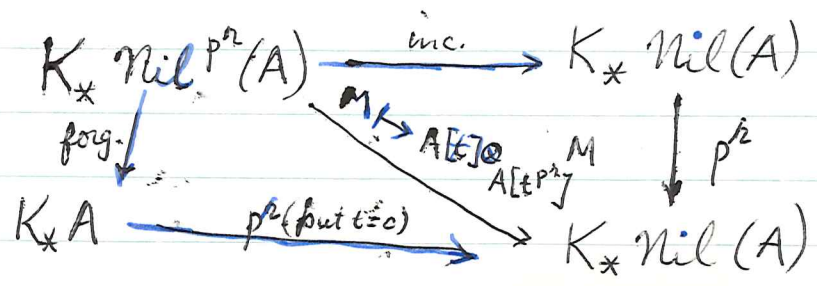
~~Now~~ Thus filtering by powers of  $I$

we get 
$$\text{gr}_j \mathbb{F}_p[t, t''] / (t'^{p^2} - t''^{p^2}) = \mathbb{F}_p[t] [t' - t''] / (t' - t'')^{p^2 - j}$$

Thus we filter:  $I^j (A[t] \otimes_{A[t^{p^2}]} M)$   $0 \leq j \leq p^2$  and the assoc.

graded  $A[t]$ -module is just  $M$  with  $t$  acting as  $\nu$ .

Thus conclude that



commutes, so done.

Question: If  $A$  is commutative, then  $A^\bullet$  acts on  $K_*(A[t])$  by  $t \mapsto \lambda t$ ,  $\lambda \in A^\bullet$ . Is this action non-trivial?

In fact this action extends to the <sup>multiplicative</sup> monoid of  $A$ , but here it is non-trivial in general; for  $\lambda=0$  gives the projection of  $K_*(A[t])$  onto  $K_*(A)$  given by  $t \mapsto 0$ .

Origin of question: To prove  $Nil_*(A)$  is torsion one might try restricting from  $A[t]$  to  $A[t^n]$  and then extending.

$$A[t] \otimes_{A[t^n]} A[t] = A[x,y]/(x^n - y^n) \quad \begin{matrix} x = t \otimes 1 \\ y = 1 \otimes t \end{matrix}$$

If  $T^n - 1$  splits into linear factors over  $A$ , then

$$x^n - y^n = \prod_{\zeta^n = 1} (x - \zeta y)$$

and so

$$\text{gr } A[t] \otimes_{A[t^n]} A[t] = \prod_{\zeta} A[t \otimes 1, 1 \otimes t] / t \otimes 1 - \zeta 1 \otimes t$$

and the effect on  $Nil$  is to replace  $(M, \nu)$  by sum of  $(M, \zeta \nu)$  as  $\zeta$  runs over roots of  $T^n - 1$ .

Question: Is it possible that for any  $\mathbb{Q}$ -algebra  $Nil$  is zero?

NO: One has:

$$K_1(A[\varepsilon]) = K_1 A \oplus A$$

where the  $A$  factors is picked up by

$$K_1(A[\varepsilon]) \xrightarrow{\det} A[\varepsilon]^\circ = A \oplus A.$$

Thus

$$\begin{aligned} \text{Nil}_0(A[\varepsilon]) &= K_1(A[\varepsilon][t]) / K_1(A[\varepsilon]) \\ &= \text{Nil}_0(A) \oplus A[t]/A \end{aligned}$$

The point is that nilpotent elts<sup>v</sup> in  $A$  contribute to units  $1+tv$  in  $A[t]$  which do not come from  $A^\circ$ .

This examples suggests:

Question: If  $l$  is a prime number invertible in  $A$ , is  $\text{Nil}_*(A)$  uniquely  $l$ -divisible?

June 11, 1973

Let  $S$  be a mult. system in the center of  $A$  consisting of non-zero divisors. Let  $P \in \mathcal{P}(S^{-1}A)$  and write

$$P \oplus Q = S^{-1}L$$

where  $L \in \mathcal{P}(A)$ . Denote  $\pi': S^{-1}L \rightarrow P$ ,  $\pi'': S^{-1}L \rightarrow Q$  the projections and put

$$L' = \pi'(L) \subset P, \quad L'' = \pi''(L) \subset Q$$

Then  $L', L''$  are finite type over  $A$ , so  $\exists s \in S$  such that

$$L \subset L' + L'' \subset s^{-1}L$$

Now suppose  $A = B[t]$ ,  $S = \{t^n\}$ , where  $B$  is regular coherent. Then

$$t^{-n}L/L \in \mathcal{P}(A/t^n A)$$

is finitely presented over  $B$ , and  $L' + L''/L$  being finitely generated over  $A$  and killed by  $t^n$  is fin. gen. over  $B$ . Since  $B$  is coherent,  $L' + L''/L$  is finitely presented over  $B$ , and since  $B$  is regular we can find a  $B[t]/t^n$  resolution

$$0 \rightarrow K \rightarrow (B[t]/t^n)^{r_m} \rightarrow \cdots \rightarrow (B[t]/t^n)^{r_0} \rightarrow L' + L''/L \rightarrow 0$$

where  $K \in \text{Nil}(B)$ . As  $\text{Nil}(B) \subset \mathcal{P}_1(B[t])$ , one has

$$L' + L''/L \in \mathcal{P}_\infty(B[t])$$

$$\Rightarrow L' + L'' \in \mathcal{P}_\infty(A)$$



But the sum  $L' + L''$  is direct, so we need

Lemma:  $P_n(A)$  is Karoubian for each  $n$ .

~~Proof: If  $R \in \text{Mod}(A)$ , then  $R \in P(A)$ .  
 $A \rightarrow \text{Hom}(R, A)$  exact and commutes with  $P_n$~~

True for  $n=0$ . Given  $L' \oplus L'' \in P_n(A)$ , then  $L' \oplus L''$  fin. gen.  $\Rightarrow L' + L''$  fin. gen.  $\Rightarrow \exists$

$$\begin{array}{ccccccc}
 & R' & & P' & & L' & \\
 0 \longrightarrow & \oplus & \longrightarrow & \oplus & \longrightarrow & \oplus & \longrightarrow 0 \\
 & R'' & & P'' & & L'' & 
 \end{array}$$

with  $P', P'' \in P(A)$ . Then get  $R' \oplus R'' \in P_{n-1}(A) \xrightarrow{\text{induction}}$   
 $R', R'' \in P_{n-1}(A) \Rightarrow L', L'' \in P_n(A)$ . done

Thus we see that  $L', L'' \in P_n(A)$  so we have proved

Prop: If  $B$  is regular ~~and~~ coherent, then  $K_0(B[t]) \rightarrow K_0(B[t, t^{-1}])$  is onto, hence  $K_0(B) \cong K_0(B[t])$

Review the argument. Given  $A$  ring and  $P \in \mathbb{P}(A[t, t^{-1}])$  write

$$P \oplus Q = S^{-1}L \quad L = A[t]^m \quad S = \{t^n\}$$

and let  $\pi': \cancel{A[t]^m} \rightarrow P$ ,  $\pi'': A[t, t^{-1}]^m \rightarrow Q$  be the two projections. Put

$$L' = \pi'(A[t]^m), \quad L'' = \pi''(A[t]^m).$$

Then have

$$A[t]^m \subset L' \oplus L'' \subset \cancel{A[t]^m} t^{-n} \cdot A[t]^m$$

for some  $n$  since  $L', L''$  are fin. gen. over  $A[t]$ .  
Now the ~~point is that~~ point is that

$$J = t^{-n} A[t]^m / L' \oplus L''$$

is a finitely presented  $A[t]$ -module ~~killed by  $t^n$~~  killed by  $t^n$ . Hence ~~is finitely presented~~  $J$  is finitely presented over  $A$ , so if we assume  $A$  regular coherent (i.e.  $\mathbb{P}_\infty(A) = \text{Mod}_{\text{fp}}(A)$ ) we have  $J \in \mathbb{P}_\infty(A)$ . This means that if we construct a resolution inductively

$$0 \rightarrow R_k \rightarrow (A[t]/t^n)^{r_k} \rightarrow \dots \rightarrow (A[t]/t^n)^{r_0} \rightarrow J \rightarrow 0$$

we eventually find that  $R_k$  is in  $\mathbb{P}(A)$ . Then by char. seq.

$$0 \rightarrow A[t] \otimes R_k \rightarrow A[t] \otimes R_k \rightarrow R_k \rightarrow 0$$

we know  $R_k \in \mathbb{P}_1(A[t])$ , so  $J \in \mathbb{P}_\infty(A[t]) \implies L' \oplus L'' \in \mathbb{P}_\infty(A[t]) \implies L' \in \mathbb{P}_\infty(A[t])$ . Since  $P = S^{-1}(L')$ , it

follows that  $cl(P)$  in  $K_0(\cancel{A} A[t, t^{-1}])$  is in the image of  $K_0(A[t])$ .

This argument seems to show that if  $S$  is a multiplicative system of central non-zero-divisors in a ring  $B$  such that ~~the image of~~

$$\mathcal{H}_S(A) = \{M \in \text{Mod}_{\text{f.p.}}(B) \mid S^{-1}M = 0\}$$

then any projective  $S^{-1}B$ -module  $P$  is of the form  $S^{-1}L$  for some  $L \in \mathcal{P}_\infty(B)$ . In fact it seems that

$$\mathcal{P}_\infty(B) \longrightarrow \mathcal{P}_\infty(S^{-1}B)$$

is onto objects. For if this be true for anything in  $\mathcal{P}_{n-1}(S^{-1}B)$ , then given  $X \in \mathcal{P}_n(S^{-1}B)$  ~~choose~~ choose  $X = P/Y$

$$\begin{array}{ccccccc} 0 & \longrightarrow & Y & \longrightarrow & P & \longrightarrow & X \longrightarrow 0 \\ & & \parallel & & \parallel & & \parallel \\ 0 & \longrightarrow & S^{-1}K & \longrightarrow & S^{-1}L & \longrightarrow & S^{-1}(L/K) \longrightarrow 0 \end{array}$$

with  $P \in \mathcal{P}(S^{-1}B)$ , so that  $Y \in \mathcal{P}_{n-1}(S^{-1}B)$ . Thus  $Y = S^{-1}K$  with  $K \in \mathcal{P}_\infty(B)$ , so  $P = S^{-1}L$ ,  $L \in \mathcal{P}_\infty(B)$  and we can assume  $K \subset L$  as  $K$  is fin. gen. Thus  $X = S^{-1}(L/K)$  where  $L/K \in \mathcal{P}_\infty(B)$ .

Probably also true that any  $K$  in  $X \in \mathcal{P}_n(S^{-1}B) \Rightarrow K$  fin. type and  $S^{-1}K = X \Rightarrow K \in \mathcal{P}_\infty(B)$ . Because we can sandwich ~~the image of~~  $L \subset K \subset S^{-1}L$ .

June 16, 1973

Let  $S$  be a cats with product  $\# \neq$  (~~ACU~~)

(H):  $S$  is a groupoid,  $\# \neq S$  faithful  $\forall S$ .

Suppose  $S$  acts on  $\mathcal{X} \neq$

- i) every arrow in  $\mathcal{X}$  is a mono.
- ii)  $\# \neq \mathcal{X}: S \rightarrow \mathcal{X}$  faithful  $\forall X \in \mathcal{X}$ .

Let  $\dots \overset{S^{-1}\mathcal{X}}{\#}$

$$p: \langle S, S \times \mathcal{X} \rangle \longrightarrow \langle S, \mathcal{X} \rangle$$

be the functor induced by <sup>the projection</sup>  $pr_2: S \times \mathcal{X} \longrightarrow \mathcal{X}$

Fix an object  $\# X_0$  of  $\langle S, \mathcal{X} \rangle$  and consider the category  $\mathcal{X}_0/p$ . An object consists of an  $(S, X)$  in  $S^{-1}\mathcal{X}$  and an arrow  $\xi: X_0 \rightarrow X$  in  $\langle S, \mathcal{X} \rangle$ , say given by

$$T \# X_0 \xrightarrow{\xi} X.$$

~~Thus an object is~~ Thus an object is

$$(S, \del{X}, T \# X_0 \xrightarrow{\xi} X)$$

and a morphism

$$(S', \del{X'}, T' \# X_0 \xrightarrow{\xi'} X')$$

is given by a  $u \in S \neq$

$$u \# S \simeq S', u \# X \simeq X'$$

~~and with  $\xi, \xi'$~~   
carrying  $\xi$  to  $\xi'$ .

To be precise: ~~an~~ an object of  $X_0 \backslash p$  consists of  $(S, X, \xi: X_0 \rightarrow X)$  and a map

$$(S, X, \xi) \longrightarrow (S', X', \xi')$$

is a map  $(S, X) \longrightarrow (S', X')$  carrying  $\xi$  to  $\xi'$ . Clearly this is cofibred over  $S^{-1}(pt)$  - composition  $(S, X, \xi) \mapsto (S, X) \mapsto S$  of two cofibred functors. The fibre over  $S$  is the eq. cat over  $X$  assoc. to functor  $X \mapsto \text{Hom}_{\langle S, X \rangle}(X_0, X)$ . But we can identify this:

$$(X, \xi) \longmapsto (T_\xi, \alpha_\xi: T_\xi \# X_0 \rightarrow X).$$

Better: Define

$$S^{-1}S \longrightarrow X_0 \backslash p$$

$$(S, T) \longmapsto ((S, T \# X_0), X_0 \rightarrow T \# X_0)$$

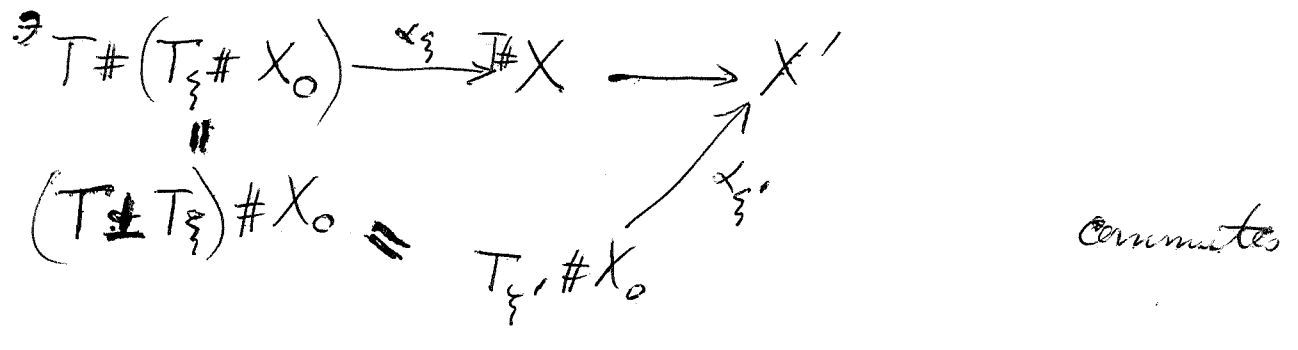
This should be a heq. In the case where  $X$  is a groupoid it is an equivalence because to give  $(S, X)$  and  $\xi: X_0 \rightarrow X$  means up to <sup>unique</sup> iso we give the  $T_\xi$  such that  $\xi$  is  $T_\xi \# X_0 \xrightarrow{\sim} X$ .

Idea: We can define a functor

$$(1) \quad X_0 \backslash p \xrightarrow{\quad} S^{-1}S$$

sending  $(S, X, \xi)$  to  $(S, T_\xi)$ . Check functoriality:

Given  $(S, X, \xi) \rightarrow (S', X', \xi')$  i.e. given  
 $T \# S \xrightarrow{\sim} S'$   
 $T \# X \rightarrow X'$

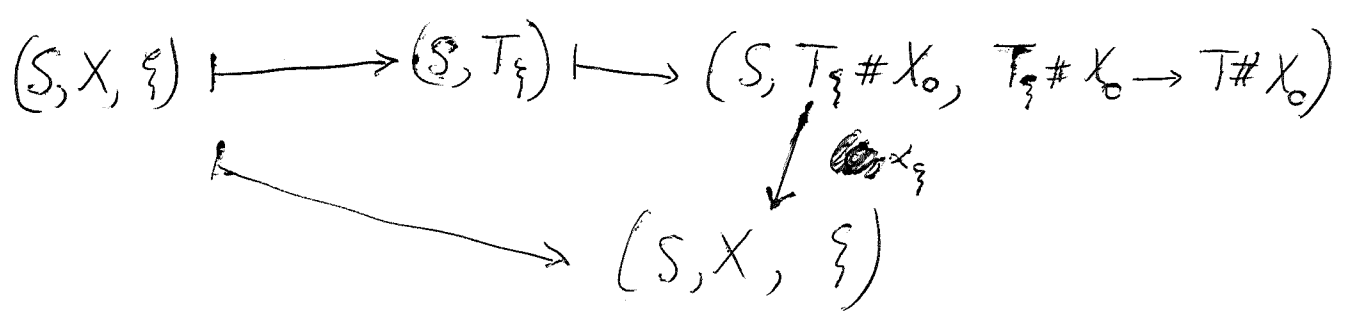


we certainly get a map  $(S, T_{\xi}) \rightarrow (S', T_{\xi'})$

On the other hand we have the functor

$$\begin{aligned}
 (2) \quad S^{-1} S &\longrightarrow X_0 \setminus p \\
 (S, T) &\longmapsto (S, T \# X_0, \text{id} : T \# X_0 \longrightarrow T \# X_0)
 \end{aligned}$$

Clearly the functor (1)(2) is isomorphic to identity. But the composition ~~(1)(2)~~ (2)(1) is homotopic to id since we have



Finally ~~if~~ if we have a map  $X_1 \xrightarrow{\beta} X_0$  in  $\langle S, X \rangle$  the following square:

$$\begin{array}{ccc}
 X_0 \setminus p & \xrightarrow{(1)} & \mathcal{S}^{-1} \mathcal{S} \\
 \downarrow & & \downarrow \text{mult by } T_\beta \\
 X_\beta \setminus p & \xrightarrow{(1)} & \mathcal{S}^{-1} \mathcal{S}
 \end{array}$$

commutes:

$$(S, X, T_\xi \# X_0 \xrightarrow{\alpha_\xi} X) \longleftrightarrow (S, T_\xi)$$

$\downarrow$

$$(S, X, T_\xi \# (T_\beta \# X_1) \rightarrow T_\xi \# X_0 \xrightarrow{\alpha_\xi} X_1) \longleftrightarrow (S, T_\xi \# T_\beta)$$

and so we see that we can apply Thm. B to conclude  $X_0 \setminus p$  is the  $h$ -fibre over  $X_0$  of the functor  $p$ .

July 11, 1973

Terminology: Category with product = symmetric monoidal category in MacLane's terminology =  $\otimes$ -cat ACU in Saavedra's terminology.

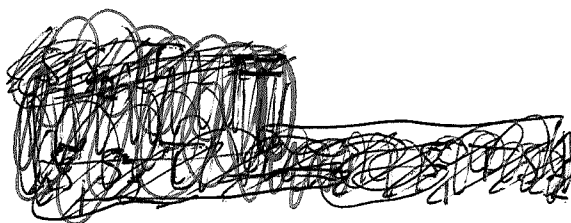
Up to equivalence we can ~~change the~~ <sup>replace</sup> such a category ~~with~~ by one in which the associativity and unity isos. are identity maps. The commutativity iso

$$\tau_{S_1 S_2}: S_1 S_2 \simeq S_2 S_1$$

is then required to satisfy:

i)  $\tau_{S_1 S_2} = \tau_{S_2 S_1}^{-1}$

ii)



$$\begin{array}{ccc} S_1 S_2 T & \simeq & T S_1 S_2 \\ \downarrow \tau & & \downarrow \tau \\ & S_1 T S_2 & \end{array}$$

iii)

$$\begin{array}{ccc} OS & \simeq & SO \\ \parallel & & \parallel \\ & S & \end{array}$$

Product-preserving functor = monoidal functor compatible with the commutativity isos.



Now suppose  $f: S \rightarrow X$  is a product-preserving functor between categories with product. Then I make  $S$  act on  $X$  in the natural way:

$$S \xrightarrow{f} X \longrightarrow \text{Hom}(X, X)$$

$\uparrow$  translation  
 left action

To show that  $\langle S, X \rangle$  and  $S^{-1}X = \langle S, S \times X \rangle$  are categories with product. Enough to do former.

Recall that  $Ob \langle S, X \rangle = Ob X$  and that a map from  $X$  to  $X'$  in  $\langle S, X \rangle$  is an eq. class of pairs  $(S, SX \rightarrow X')$ . Define a product functor

$$\langle S, X \rangle \times \langle S, X \rangle \longrightarrow \langle S, X \rangle$$

$$(X_1, X_2) \longmapsto X_1 X_2$$

$$(cl(S_1, S_1 X_1 \rightarrow X'_1), cl(S_2, S_2 X_2 \rightarrow X'_2)) \longmapsto \text{class of}$$

$$(S_1 S_2, (S_1 S_2)(X_1 X_2) \xrightarrow{\cong} S_1 X_1 S_2 X_2 \rightarrow X'_1 X'_2)$$

Here to save writing we put  $SX$  for  $f(S) \cdot X$  and ~~the iso~~ the iso is

$$\begin{aligned}
 f(S_1 S_2)(X_1 X_2) &\cong (f S_1 \cdot f S_2)(X_1 X_2) && \text{product iso} \\
 &= f S_1 (f S_2 (X_1 X_2)) && \text{assoc} \\
 &= f S_1 ((f S_2 \cdot X_1) X_2) && \text{assoc.}
 \end{aligned}$$

$$\begin{aligned}
&\simeq f_{S_1}((X_1; f_{S_2})X_2) && \text{Comm} \\
&= f_{S_1}(X_1, (f_{S_2}X_2)) && \text{assoc} \\
&= (f_{S_1} \cdot X_1)(f_{S_2} \cdot X_2). && \text{assoc.}
\end{aligned}$$

If  $S, X, f$  are struct it is simply the ~~map~~ map  
 $f(S_1 S_2) X_1 X_2 = f_{S_1} f_{S_2} X_1 X_2 \simeq f_{S_1} X_1' \cdot f_{S_2} X_2$

induced by <sup>the</sup> commutativity isom of  $X$ .

~~It~~ It is clear this procedure above does carry equivalence classes to equiv. classes. We must check it is compatible with composition. So suppose given

$$\begin{array}{ll}
S_1 X_1 \rightarrow X_1' & S_2 X_2 \rightarrow X_2' \\
S_1' X_1' \rightarrow X_1'' & S_2' X_2' \rightarrow X_2''
\end{array}$$

Composing ~~maps~~ <sup>we</sup> get

~~the~~

$$S_1' S_1 X_1 \rightarrow S_1' X_1' \rightarrow X_1'', \quad S_2' S_2 X_2 \rightarrow S_2' X_2' \rightarrow X_2''$$

which goes to

$$S_1' S_1 S_2' S_2 X_1 X_2 \simeq S_1' S_1 X_1 S_2' S_2 X_2 \rightarrow X_1'' X_2''$$

On the other hand composing

$$S_1 S_2 X_1 X_2 \simeq S_1 X_1 S_2 X_2 \longrightarrow X'_1 X'_2$$

$$S'_1 S'_2 X'_1 X'_2 \simeq S'_1 X'_1 S'_2 X'_2 \longrightarrow X''_1 X''_2$$

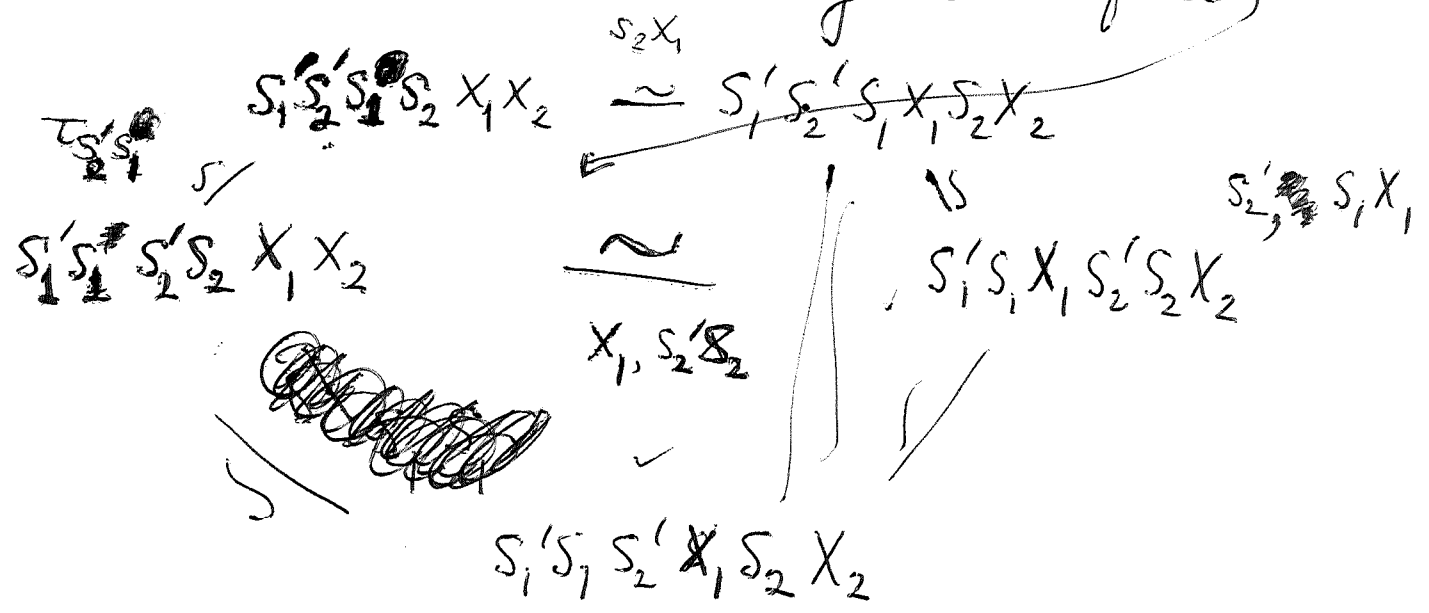
gives

$$S'_1 S'_2 S_1 S_2 X_1 X_2 \simeq \underbrace{S'_1 S'_2 S_1 X_1 S_2 X_2}_{\text{IS}} - \underbrace{S'_1 S'_2 X'_1 X'_2}_{\text{IS}}$$

$$S'_1 S_1 X_1 S'_2 S_2 X_2 \longrightarrow S'_1 X'_1 S'_2 X'_2$$

↓  
X''\_1 X''\_2

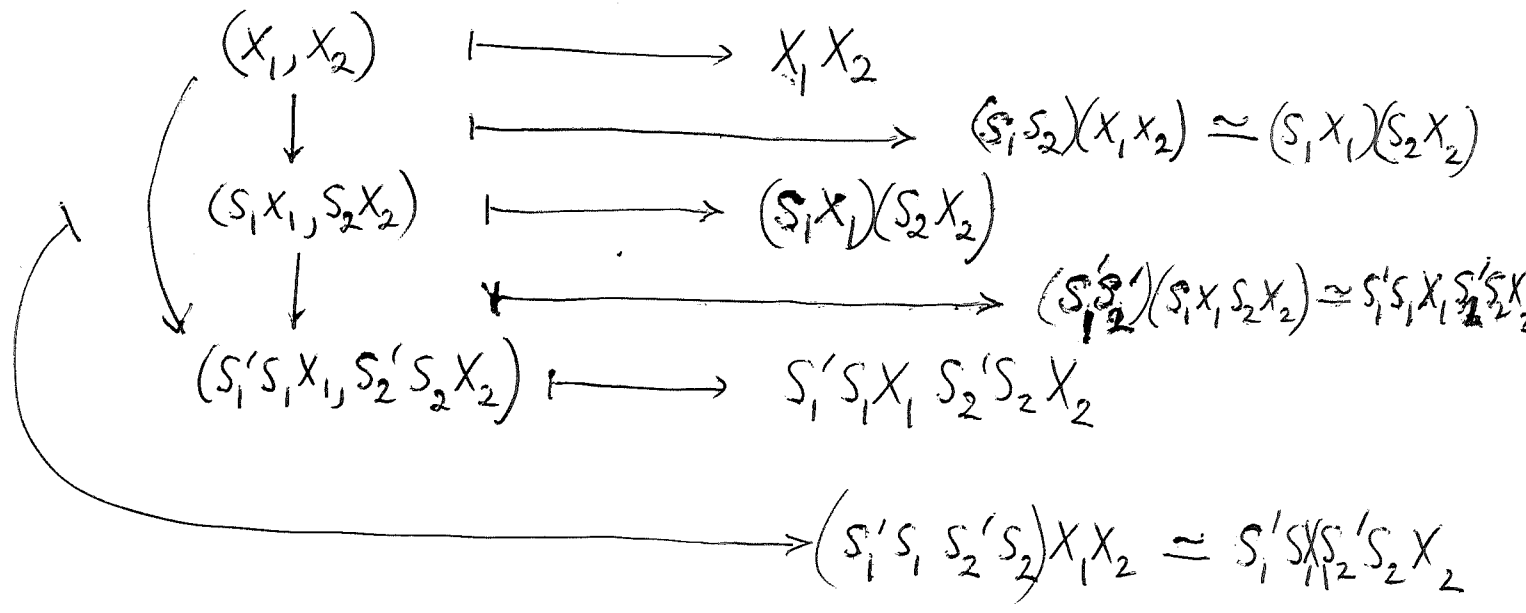
So to see these two results are the same we must establish commutativity in the square)



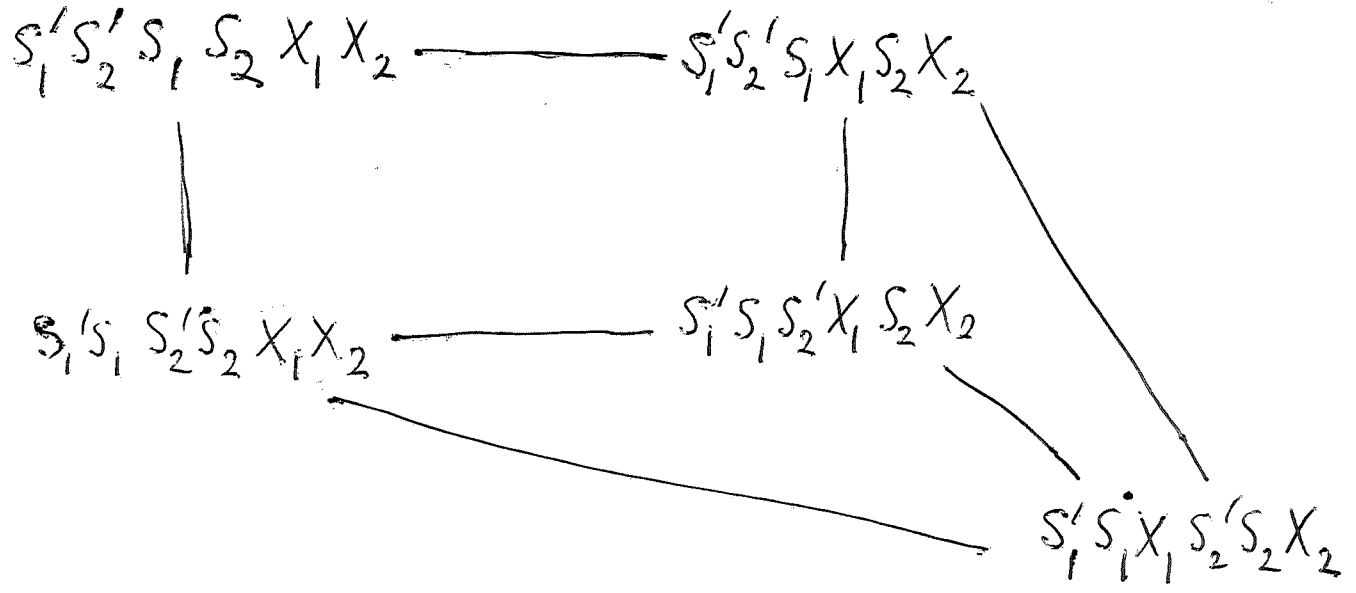
which is clear.

What matters is that the composition of

~~$(X_1, X_2)$~~



Thus what we really have to show is that the outside square in



commutes. Clear as the triangles commute by hyp.

June 30, 1973.

monoidal categories.

Definition of monoidal category (terminology of MacLane; it is called  $\otimes$ -cat  $\mathcal{A}$  in Saavedra). A monoidal category ~~is a cat  $\mathcal{S}$  equipped with the following extra~~ ~~structure~~ consists of the following data:

- i) ~~a category  $\mathcal{S}$~~
- ii) a functor  $\perp: \mathcal{S} \times \mathcal{S} \rightarrow \mathcal{S}$ , ~~denoted~~ denoted  $(S_1, S_2) \mapsto S_1 \perp S_2$  or simply  $(S_1, S_2) \mapsto S_1 S_2$ , and called the product functor of the monoidal category
- iii) an isomorphism

$$\alpha: (S_1 S_2) S_3 \xrightarrow{\sim} S_1 (S_2 S_3)$$

of functors from  $\mathcal{S} \times \mathcal{S} \times \mathcal{S}$  to  $\mathcal{S}$  called the associativity isom of ~~the monoidal cat.~~ the monoidal cat.

- iii) an object  $0$  of  $\mathcal{S}$ , called the unit of  $\mathcal{S}$  and two isomorphisms

$$0S \xrightarrow{\sim} S, \quad S0 \xrightarrow{\sim} S$$

of functors from  $\mathcal{S}$  to  $\mathcal{S}$ , called the left and right unit isom, respectively, which coincide for  $S=0$ .

These data are subject to the following conditions:

- 1) The associativity pentagon

$$\begin{array}{ccc} ((S_1 S_2) S_3) S_4 & \xrightarrow{\sim} & (S_1 (S_2 S_3)) S_4 \\ \downarrow & & \downarrow \\ (S_1 S_2) (S_3 S_4) & & S_1 (S_2 S_3) S_4 \end{array}$$

$$(S_1 S_2) (S_3 S_4)$$

$\downarrow$

$$S_1 (S_2 (S_3 S_4)) \xrightarrow{\sim} S_1 ((S_2 S_3) S_4)$$

should commute

~~Alternative language: A monoidal structure on a cat.  $\mathcal{C}$  consists of the following data~~

Two possible  $00 \approx 00$  coincide

2) The triangles

$$(0 S_1) S_2 \cong 0(S_1 S_2)$$

$$\searrow \quad \swarrow$$

$$S_1 S_2$$

$$(S_1 0) S_2 \cong S_1(S_2 0)$$

$$\searrow \quad \swarrow$$

$$S_1 S_2$$

$$(S_1 S_2) 0 \cong S_1(S_2 0)$$

$$\searrow \quad \swarrow$$

$$S_1 S_2$$

should commute. (According to MacLane the middle ones & assoc. pentagon  $\Rightarrow$  other two).

~~Example: Let  $\mathcal{C}$  be a monoidal object in  $\mathcal{C}at$  (a 2-category with a unique object). This gives a monoidal cat  $\mathcal{C}$  is a category equipped with a monoidal structure.~~

Example: Call a monoidal cat strict if the associativity and unit isos are identity maps. Clearly strict monoidal cats are the same thing as monoid-objects in Cat (which can also be identified with 2-categories with a unique object), for example Hom( $X, X$ ).

Definition of monoidal functor: Let  $S, T$  be monoidal categories, ~~By a monoidal functor  $F: S \rightarrow T$  we mean a functor  $S \rightarrow T$  equipped with the following extra structure:~~ By a monoidal functor  $F: S \rightarrow T$  we mean a functor  $S \rightarrow T$  (from  $S$  to  $T$ ) equipped with the following

extra structure:  
i) an isomorphism  $F(S_1 S_2) \cong F(S_1) F(S_2)$  of functors from  $S \times S$  to  $T$  called the product isom of  $F$ .

ii) an isom  $F(0) \simeq 0$  called the unit isom.  
of  $F$ .

These data are subject to the conditions:

1) (compatibility with associativity data):

$$\begin{array}{ccc} F((S_1, S_2), S_3) & \simeq & F(S_1, S_2) F(S_3) \simeq [F(S_1) F(S_2)] F(S_3) \\ \downarrow \cong & & \downarrow \cong \\ F(S_1, (S_2, S_3)) & \simeq & F(S_1) F(S_2, S_3) \simeq F(S_1) [F(S_2) F(S_3)] \end{array}$$

2) (compatibility with unit data)

$$\begin{array}{ccc} F(0, S) \simeq F(0) F(S) & & F(S, 0) \simeq F(S) F(0) \\ \downarrow \cong & \downarrow \cong & \downarrow \cong \quad \downarrow \cong \\ F(S) \simeq 0 F(S) & & F(S) \simeq F(S) 0 \end{array}$$

Definition of natural transf of monoidal functors:  
Let  $F, G: \mathcal{S} \rightarrow \mathcal{T}$  be two monoidal functors.  
By a <sup>monoidal</sup> natural transf of monoidal functors  $u: F \rightarrow G$   
we mean a natural transf of functors such that

$$\begin{array}{ccc} F(S_1, S_2) \longrightarrow G(S_1, S_2) & & F(0) \longrightarrow G(0) \\ \downarrow \cong & \downarrow \cong & \searrow \quad \swarrow \cong \\ F(S_1) F(S_2) \longrightarrow G(S_1) G(S_2) & & 0 \end{array}$$

commutes.

Composition of monoidal functors. Given monoidal functors

$$\begin{aligned}
 F: \mathcal{S} &\longrightarrow \mathcal{T} & S &\longmapsto FS, & F(S_1, S_2) &\simeq F(S_1)F(S_2) \\
 G: \mathcal{T} &\longrightarrow \mathcal{U} & T &\longmapsto GT, & G(T_1, T_2) &\simeq G(T_1)G(T_2)
 \end{aligned}$$

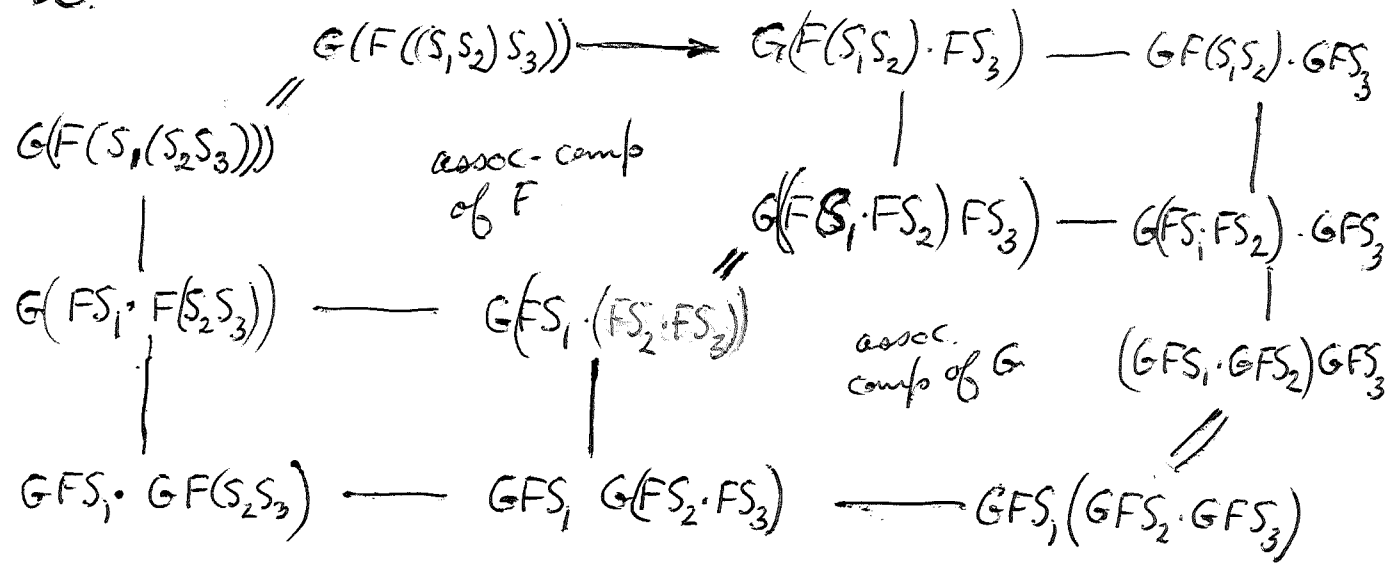
let  $GF$  be equipped with the product via

$$(GF)(S_1, S_2) = G(F(S_1, S_2)) \simeq G(F(S_1)F(S_2)) \simeq GF(S_1)GF(S_2)$$

and the unit is

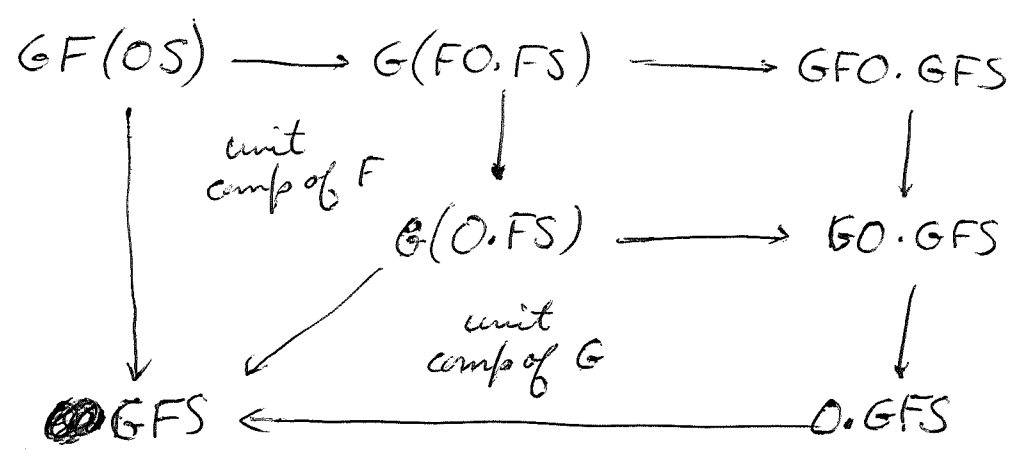
$$(GF)(0) = G(F(0)) \simeq G(0) \simeq 0.$$

Claim then  $GF$  is a monoidal functor from  $\mathcal{S}$  to  $\mathcal{U}$ .



The above <sup>comm</sup> diagram shows  $GF$  is compatible with associativity data





The above commutative diagram shows that  $GF$  is compatible with the unit data.

Put  $\text{Hom}_{\text{monoidal}}^{\downarrow}(S, T)$  for the category of monoidal functors from  $S$  to  $T$  and natural transfs. Clearly the composition just defined gives rise to a functor

$$\begin{array}{ccc}
 \underline{\text{Hom}}^{\downarrow}(S, T) \times \underline{\text{Hom}}^{\downarrow}(T, U) & \longrightarrow & \underline{\text{Hom}}^{\downarrow}(S, U) \\
 (F, G) & \longmapsto & GF
 \end{array}$$

Evident transitivity requirements hold, so it is clear that we have established

Assertion: Monoidal categories, monoidal functors, and their <sup>monoidal</sup> natural transformations form a 2-category.

So we now ~~define an equivalence~~ say that a monoidal functor  $F: S \rightarrow T$  is an equivalence of monoidal categories if  $\exists$  a monoidal functor  $G: T \rightarrow S$  and monoidal isos.

$$GF \cong id_S, \quad FG \cong id_T.$$

Lemma: A monoidal functor  $F: \mathcal{S} \rightarrow \mathcal{T}$  is an equivalence of monoidal categories iff it is an equivalence of categories.

Proof: If  $G$  ~~is~~ <sup>is</sup> a functor quasi-inverse to  $F$ , then ~~we will have~~

~~$G(T_1 \cdot T_2) \simeq G(T_1) \cdot G(T_2)$~~   
 ~~$G(F(T_1) \cdot F(T_2)) \simeq G(F(T_1)) \cdot G(F(T_2))$~~   
 ~~$G(F(G(T_1) \cdot G(T_2))) \simeq G(F(G(T_1))) \cdot G(F(G(T_2)))$~~

to the isomorphism in  $\mathcal{T}$

$$F(GT_1 \cdot GT_2) \simeq FGT_1 \cdot FGT_2 \simeq T_1 \cdot T_2$$

corresponds an isomorphism

$$(1) \quad GT_1 \cdot GT_2 \simeq G(T_1 T_2)$$

in  $\mathcal{S}$ . And similarly to

$$F(GO) \simeq O \simeq FO$$

corresponds (2):  $GO \simeq O$ . Better: (1) <sup>(2)</sup> use the unique maps such that

$F(GT_1, GT_2)$	$\xrightarrow{F(1)}$	$FG(T_1, T_2)$		$F(GO)$	$\xrightarrow{F(2)}$	$FO$
$\downarrow \text{product iso of } F$		$\downarrow \text{ }$		$\downarrow$		$\downarrow \text{ unit iso of } F$
$FGT_1 \cdot FGT_2$		$\simeq$	$T_1 \cdot T_2$			$O$

commute. To show ~~the~~ that (1) (2) constitutes a monoidal functor structure on  $G$  we must prove

that certain diagrams constructed from these maps commute. But because  $F$  transports (1) (2) into the product iso of  $F$  and unit iso of  $F$  ~~and~~ the diagrams in question commute, since they already do so for  $F$ .

~~But we use the equivalence  $\mathcal{C}$  to transport~~

Definition: An <sup>(left)</sup> action of a monoidal category  $S$  on a category  $\mathcal{X}$  consists of  
 i) a functor  $\# : S \times \mathcal{X} \rightarrow \mathcal{X}$  denoted  $(S, X) \mapsto S\#X$ ,  
 or simply  $(S, X) \mapsto SX$ .  
 ii) an isom

$$(S_1 S_2) X \xrightarrow{\sim} S_1 (S_2 X)$$

of functors  $S \times S \times \mathcal{X} \rightarrow \mathcal{X}$ , called the assoc isom of the action

iii) an isom

$$0X \xrightarrow{\sim} X$$

of functors  $\mathcal{X} \rightarrow \mathcal{X}$  called the unit isom of the action.

These data are supposed to satisfy

1) ~~Associativity~~

$$(S_1 S_2 S_3) X \xrightarrow{\sim} (S_1 S_2) (S_3 X) \xrightarrow{\sim} S_1 (S_2 (S_3 X))$$

$$(S_1 (S_2 S_3)) X \xrightarrow{\sim} S_1 ((S_2 S_3) X)$$

Commutates

$$\begin{array}{ccc}
 2) & (S0)X \simeq S(0X) & (0S)X \simeq 0(SX) \\
 & \searrow \quad \swarrow & \searrow \quad \swarrow \\
 & SX & SX
 \end{array}$$

commute.

Example:  $\exists$  canonical action of  $S$  on itself with ~~product~~ <sup>action</sup> functor  $(S, X) \mapsto SX$  the product functor of  ~~$S$~~   $S$ , with assoc. iso. the assoc iso of  $S$ , and with unit isom. the <sup>left</sup> unit isom  $0X = X$  of  $S$ .

~~Alternative definition: Given an action of  $S$  on  $X$  and  $\eta$  a monoidal functor  $S \rightarrow \text{Hom}(X, X)$  follows.~~

Alternative definition: An action of  $S$  on  $X$  may be identified with a monoidal functor

$$S \rightarrow \text{Hom}(X, X).$$

In effect the <sup>action</sup> functor  $S \times X \rightarrow X$ ,  $(S, X) \mapsto SX$  may be identified with a functor  $S \rightarrow \text{Hom}(X, X)$ ,  $S \mapsto (X \mapsto SX)$ . The assoc. isom. for the ~~action~~ <sup>action</sup> ~~corresponds to~~ <sup>corresponds to</sup> product iso for the functor, and similarly for the unit iso. Conditions (1) and (2) of definition of monoidal functor and action coincide

Let  $\mathcal{X}, \mathcal{Y}$  be cats with  $\text{Set}_{\mathcal{X}}$   
Definition of action-preserving functor: An  
 action-preserving functor  $F: \mathcal{X} \rightarrow \mathcal{Y}$  is a functor  
 equipped with an isomorphism

$$S F(X) \xrightarrow{\sim} F(SX) \quad \text{action-preserving iso}$$

of functors  $S \times \mathcal{X} \rightarrow \mathcal{Y}$  satisfying

1) (Compat with assoc. <sup>iso</sup> of the action)

$$\begin{array}{ccc} (S_1 S_2) F(X) & \xrightarrow{\sim} & F((S_1 S_2) X) \\ \parallel & & \parallel \\ S_1 (S_2 F(X)) & \xrightarrow{\sim} & S_1 F(S_2 X) \xrightarrow{\sim} F(S_1 (S_2 X)) \end{array}$$

2) (compatible with unit iso. for the action)

$$\begin{array}{ccc} 0 F(X) & \xrightarrow{\sim} & F(0X) \\ \parallel & & \parallel \\ & & F(X) \end{array}$$

Given action-preserving functors  $F, G: \mathcal{X} \rightarrow \mathcal{Y}$   
 a morphism  $F \rightarrow G$  of action-preserving functors  
 is a natural transf compatible with the action-preserving  
 isom:

$$\begin{array}{ccc} S F(X) \xrightarrow{\sim} F(SX) & & \\ \downarrow & & \downarrow \\ S G(X) \xrightarrow{\sim} G(SX) & & \end{array}$$

Category  $\text{Hom}^{\text{act}}(\mathcal{X}, \mathcal{Y})$  of action-preserving functors

Composition: Given action-preserving functors

$$X \xrightarrow{F} Y \xrightarrow{G} Z$$

we equip  $GF$  with the isom

\*  $S(GF(X)) \simeq G(SF(X)) \simeq S(GF(X)).$

Verify that in this way  $GF$  is an action-preserving functor:

$$\begin{array}{ccccc}
 (S_1 S_2)GF X & \simeq & G(S_1 S_2)FX & \simeq & GF(S_1 S_2)X \\
 \parallel & \text{because } G \text{ is} & \parallel & & \parallel \\
 & \text{action-preserving} & & & \text{because} \\
 S_1 S_2 GF X & \simeq & S_1 G S_2 FX & \simeq & G S_1 S_2 FX & \text{F is action-preserving} \\
 & \parallel & \text{triv.} & \parallel & \\
 S_1 G F S_2 X & \simeq & G S_1 F S_2 X & \simeq & G F S_1 S_2 X
 \end{array}$$

The above diagram shows  $GF$  <sup>with \*</sup> comp. with assoc iso.

$$\begin{array}{ccccc}
 OGF X & \simeq & GOF X & \simeq & GFO X \\
 \searrow & \text{because } G & \downarrow & \text{because} & \swarrow \\
 & \text{is act.-pres.} & GF X & \text{F is action} & \\
 & & & \text{preserving} & 
 \end{array}$$

The above shows  $(GF, *)$  is compat. with unit isos.

Clearly composition defines a functor

$$\begin{array}{ccc} \underline{\text{Hom}}^{\mathcal{S}}(\mathcal{X}, \mathcal{Y}) \times \underline{\text{Hom}}^{\mathcal{S}}(\mathcal{Y}, \mathcal{Z}) & \longrightarrow & \underline{\text{Hom}}^{\mathcal{S}}(\mathcal{X}, \mathcal{Z}) \\ (F, G) & \longmapsto & GF \end{array}$$

which is associative with identities. Thus have verified

Assertion: Categories with  $\mathcal{S}$ -action, ~~action-preserving~~ action-preserving functors, and ~~morphisms~~ morphisms of action-preserving functors form a 2-category.

In particular  $\underline{\text{Hom}}^{\mathcal{S}}(\mathcal{X}, \mathcal{X})$  is a strict monoidal category.

It should be clear what a right action of a monoidal cat.  $\mathcal{S}$  on a cat.  $\mathcal{X}$  is. It is a left action of  $\mathcal{S}^*$  on  $\mathcal{X}$ , where  $\mathcal{S}^*$  denotes the category  $\mathcal{S}$  equipped with the opposite monoidal structure, i.e. in which the functor  $\perp: \mathcal{S}^* \times \mathcal{S}^* \rightarrow \mathcal{S}^*$  is  $(S, T) \mapsto TS$ .

Example:  $\mathcal{S}^*$  acts on  $\mathcal{S}$  ~~with~~ with product

$$\mathcal{S}^* \times \mathcal{S} \longrightarrow \mathcal{S}, \quad (S, X) \longmapsto XS$$

and assoc. + unity isos. obtained from those of  $\mathcal{S}$  in the evident way.

~~Proposition 1.1~~

We propose now to define a monoidal functor

$$(*) \quad \lambda: \mathcal{S} \longrightarrow \underline{\text{Hom}}^{\mathcal{S}^*}(\mathcal{S}, \mathcal{S})$$

where  $\mathcal{S}^*$  acts on  $\mathcal{S}$  in the way just described. Given an object  $S$  one ~~equips~~ <sup>equips</sup> the functor  $X \mapsto SX$  from  $\mathcal{S}$  to itself together with the isomorphism

$$(SX)T \simeq S(XT)$$

of functors from  $\mathcal{S} \times \mathcal{S}^*$  to  $\mathcal{S}$  given by the assoc. isom of  $\mathcal{S}$ . In this way we obtain an  $\mathcal{S}^*$ -action-preserving functor from  $\mathcal{S}$  to  $\mathcal{S}$ .

$$\lambda S = (X \mapsto SX, (SX)T \simeq S(XT)).$$

(The ~~two~~ <sup>two</sup> axioms for ~~action-preserving~~ <sup>action-preserving</sup> result from the axioms for a monoidal cat). In this way we obtain a functor  $(*)$ .

To make  $(*)$  a monoidal functor we equip  $(*)$  with the isom

$$(**) \quad \lambda(S_1 \otimes S_2) \xrightarrow{\sim} \lambda S_1 \cdot \lambda S_2$$

given by the assoc. isom.

$$(S_1 \otimes S_2)X \xrightarrow{\sim} S_1(S_2X)$$

It is necessary to check this is an isom of  $\mathcal{S}^*$ -action-preserving functors. Thus we must show commutativity of



$$((S_1 S_2) X) T \simeq (S_1 S_2)(X T)$$

$$\downarrow$$

$$\downarrow$$

$$(S_1(S_2 X)) T \simeq S_1((S_2 X) T) \simeq S_1(S_2(X T))$$

(bottom row is the action com for  $\lambda_{S_1} \cdot \lambda_{S_2}$ ), which results from the pentagon axiom.

We also must equip  $\lambda$  with

$$(**)' \quad \lambda 0 \simeq \text{id}$$

$$\text{given by} \quad 0 X \simeq X.$$

To see this is an isomorphism of  $S^*$ -action-preserving functors, we need to check commutativity of

$$\textcircled{0} (0 X) T \simeq 0(X T) \textcircled{0}$$

$$\downarrow$$

$$\downarrow$$

$$X T = X T$$

which is clear.

Now the verification that  $(**)$  and  $(**)'$  make  $\lambda$  into a monoidal functor requires checking the comm. of certain diagrams in  $\underline{\text{Hom}}^{S^*}(S, S)$  and this can be verified objectwise. So it's all clear now that  $\lambda$  is a monoidal functor.

Proposition:  $\lambda : S \rightarrow \underline{\text{Hom}}^{S^*}(S, S)$  is an equivalence of monoidal categories. In particular any monoidal category is equivalent (as monoidal cat.) to a strict monoidal category.

Proof. By the lemma it suffices to show  $\lambda$  is an equivalence of categories. Define a functor

$$\mu: \text{Hom}^{\mathcal{S}^*}(\mathcal{S}, \mathcal{S}) \longrightarrow \mathcal{S}$$

$$\mu(X \mapsto FX, (FX)T \simeq F(XT)) = FO.$$

Then

$$\mu\lambda(S) = \mu(X \mapsto SX, (SX)T \simeq S(XT))$$

$$= SO$$

is isomorphic to the id functor of  $\mathcal{S}$ . Also

$$\lambda\mu(F) = (X \mapsto (FO)X, (FO \cdot X)T \simeq \overset{FO(XT)}{\cancel{(FO)XT}})$$

Because  $F$  is  $\mathcal{S}^*$  action-preserving we have a canonical isom

$$(FO)X \simeq F(OX) \simeq FX$$

of functors from  $\mathcal{S}$  to  $\mathcal{S}$ . To show this is a morphism of action-preserving functors we have to verify:

$$((FO)X)T \simeq (FO)(XT)$$

$$\downarrow \quad \quad \quad \downarrow$$

$$F(OX)T \simeq F(OXT) \simeq F(O(XT))$$

$$\downarrow \quad \quad \quad \downarrow \quad \quad \quad \downarrow$$

$$(FX)T \simeq F(XT)$$

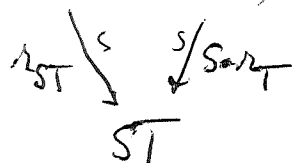
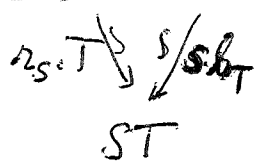
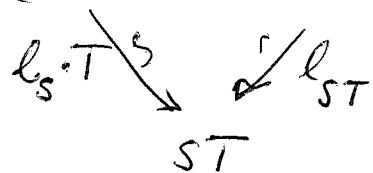
So done.

What must be true therefore is that

$$(OS)T \simeq O(ST)$$

$$(SO)T \simeq S(OT)$$

$$(ST)O \simeq S(TO)$$

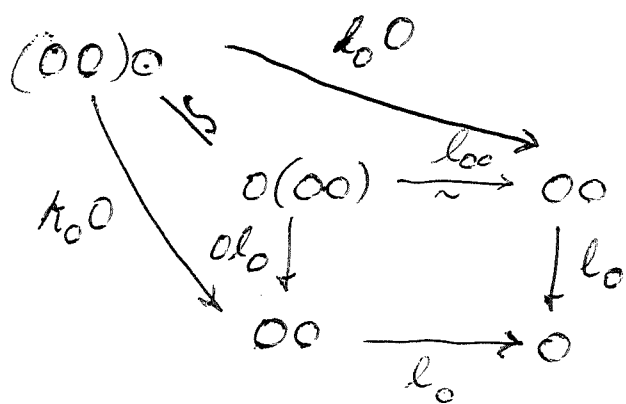


must imply  
direct proof:

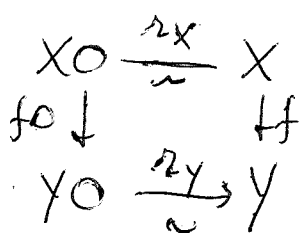
$$r_0 = l_0 : 00 \rightarrow 0$$

Here is a

start with



Since  $l_0$  is an isom  $\Rightarrow r_0 O = l_0 O$ . But because of



hence  $f0 = g0 \Leftrightarrow f = g$

we have  $r_0 O = l_0 O \Rightarrow r_0 = l_0$ .

~~Below shows that any monoidal category~~

Above shows that any monoidal category is equivalent to a strict one. Given a word  $w$  in the free <sup>non-</sup>associative ~~by~~ set with one generator

e.g.  $(x(xx))x)x$

it determines a functor  $S^n \rightarrow S$ , where  $n = \text{length of } w$ . Using the associativity isom. one gets ~~the~~ the notion of an elementary transition between words, and composing these one gets various isomorphisms between the functors  $S^n \rightarrow S$  associated to the different words. The coherence theorem states that any two such isos between the functors assoc. to two given words coincide. To prove this one reduces immediately to ~~the~~ the case where  $S$  is strict, whence all words of length  $n$  lead to same functor  $S^n \rightarrow S$ , and all <sup>transition</sup> isos. between these are the identity.

---

Let  $S$  be a monoidal category, let

$$\tilde{S} = \underline{\text{Hom}}^{S^*}(S, S)$$

so that we have the equivalence  $S \xrightarrow{\sim} \tilde{S}$  as object. We define a new category  $\tilde{S}$  as follows,  $\text{Ob}(\tilde{S}) = \text{free monoid generated by } \text{Ob}(S)$ . Since  $\tilde{S}$  is strict, the canon map  $S \rightarrow \tilde{S}$  induces a map of monoids  $\text{Ob}(\tilde{S}) \rightarrow \text{Ob}(\tilde{S})$ .

In other words an element of  $\text{Ob}(\tilde{\mathcal{S}})$  is a word  $S_1 \dots S_n$  of objects of  $\mathcal{S}$  ( $n \geq 0$ ) and this word gets sent to the functor

$$X \mapsto S_1(S_2 \dots (S_n X) \dots)$$

with its evident  $\mathcal{S}^*$ -action-preserving structure. Now we define morphisms in  $\tilde{\mathcal{S}}$  so that

$$\tilde{\mathcal{S}} \longrightarrow \bar{\mathcal{S}}$$

$$S_1 \dots S_n \mapsto (X \mapsto S_1(S_2 \dots (S_n X) \dots))$$

is an equivalence of categories. Define product  $\tilde{\mathcal{S}} \times \tilde{\mathcal{S}} \longrightarrow \tilde{\mathcal{S}}$  functor by sending

$$f: S_1 \dots S_n \longrightarrow S'_1 \dots S'_m \quad \text{i.e.}$$

$$S_1 S_2 \dots S_n X \longrightarrow S'_1 \dots S'_m X$$

$$g: T_1 \dots T_p \longrightarrow T'_1 \dots T'_q$$

$$T_1 \dots T_p X \longrightarrow T'_1 \dots T'_q X$$

into

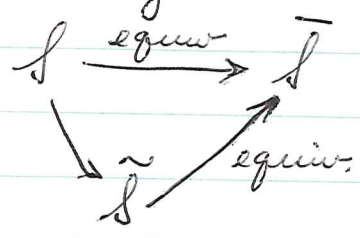
$$f \cdot g: S_1 \dots S_n T_1 \dots T_p \longrightarrow S'_1 \dots S'_m T'_1 \dots T'_q$$

the composition of the natural transformations. (Lemma here is that if  $M$  is a monoid,  $\mathcal{S}$  is a <sup>strict</sup> monoidal category, and if  $M \longrightarrow \mathcal{A}(\mathcal{S})$  is given, then the category  $\mathcal{M}$  having  $M$  for objects and over  $\mathcal{S} \ni M \longrightarrow \mathcal{S}$  is fully faithful is a monoidal cat. In fact we have

$$\mathcal{M} = D(M) \times_{D(\text{Ob } \mathcal{S})} \tilde{\mathcal{S}}$$

where  $D(M)$  is category with objects =  $M$ , arrows =  $M \times M$ , etc

Thus we see  $\tilde{\mathcal{S}}$  is a strict monoidal cat. and the functor  $\mathcal{S} \rightarrow \tilde{\mathcal{S}}$  is an equivalence of monoidal categories. This is because



where equiv means an equivalence of monoidal categories.

Now let  $f: \mathcal{S} \rightarrow \mathcal{T}$  be a monoidal functor. If  $\mathcal{T}$  is strict, then we get a monoid map

$$f: \text{Ob}(\tilde{\mathcal{S}}) \rightarrow \text{Ob}(\mathcal{T})$$

extending the given map  $f: \text{Ob}(\mathcal{S}) \rightarrow \text{Ob}(\mathcal{T})$ . Now if  $S_1 \cdots S_n$  is an object of  $\mathcal{S}$ , we know it is canonically isomorphic to  $S_1 + (S_2 + \cdots + S_n)$  of  $\tilde{\mathcal{S}}$ , hence  $\mathcal{T}$  canonically

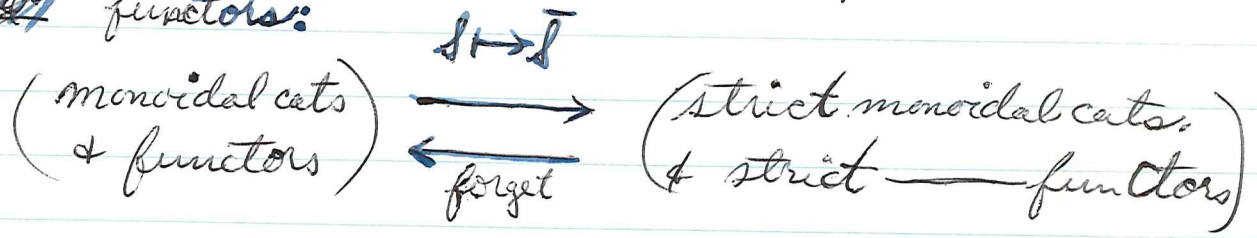
$$f(S_1 + (\cdots + S_n)) \simeq f(S_1) \cdots f(S_n) = f(S_1 \cdots S_n) \quad \because$$

as  $f$  is monoidal. Thus since  $\mathcal{S} \rightarrow \tilde{\mathcal{S}}$  is an equivalence, we can extend  $f$  to a functor  $f: \tilde{\mathcal{S}} \rightarrow \mathcal{T}$  such that

$$\begin{array}{ccc}
 f(S_1 \cdots S_n) & = & f(S_1) \cdots f(S_n) \\
 \downarrow & & \downarrow \\
 & & f(S_1 + (S_2 + \cdots + S_n))
 \end{array}$$

Thus it is fairly clear that  $f: \tilde{\mathcal{S}} \rightarrow \mathcal{T}$  has to be a strict

monoidal functors. ~~In fact, it seems we have adjoint functors  $\mathcal{S} \rightarrow \tilde{\mathcal{S}}$  is a monoidal map to a strict monoidal cat.~~



Improvement: Given  $f: \mathcal{S} \rightarrow \mathcal{T}$  monoidal with  $\mathcal{T}$  strict I wish to show  $\exists$  unique  $g: \tilde{\mathcal{S}} \rightarrow \mathcal{T}$  strict monoidal such that  $g|_{\mathcal{S}} = f$

Uniqueness of  $g$ : Because  $\mathcal{S} \rightarrow \tilde{\mathcal{S}}$  is an equivalence, given  $g_1, g_2$  there is a unique isom.  $\theta: g_1 \xrightarrow{\sim} g_2$  inducing the identity on  $\mathcal{S}$ . To show  $\theta$  is the identity. Given  $S_1 \dots S_n \in \mathcal{S}$

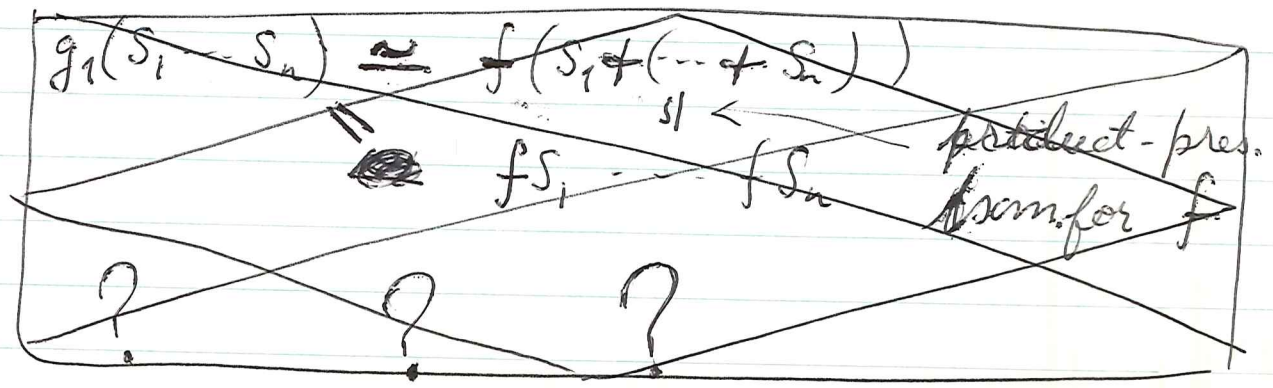
$$\begin{aligned}
 g_i(S_1 \dots S_n) &= g_i(S_1) \dots g_i(S_n) & i=1,2 \\
 &= fS_1 \dots fS_n
 \end{aligned}$$

But in  $\tilde{\mathcal{S}}$  there is an isom canonical

$$S_1 \dots S_n \cong S_1 \ast (S_2 \ast \dots \ast S_n) \dots$$

product iso for  $\mathcal{S} \rightarrow \tilde{\mathcal{S}}$

with  $\ast =$  product in  $\mathcal{S}$ . Thus



$$\begin{array}{ccc}
 & g_1(\text{product isom}) & \\
 & \downarrow \text{for } S \rightarrow \tilde{S} & \\
 g_1(S_1 \cdots S_n) & \xrightarrow{\sim} & g_1(S_1 + (\cdots + S_n)) = f(S_1 + (\cdots + S_n) \cdots) \\
 \text{product iso} \parallel & & \parallel \text{product iso for } f \\
 \text{for } g & g_1 S_1 \cdots g_1 S_n & f S_1 f S_2 \cdots f S_n
 \end{array}$$

To avoid confusion introduce  $c: S \rightarrow \tilde{S}$ . Then

$$c S_1 \cdots c S_n \xrightarrow{\sim} c(S_1 + (\cdots + S_n)) \quad \text{product iso for } c: S \rightarrow \tilde{S}$$

$$\begin{array}{ccc}
 \text{product iso. for } g \rightarrow \parallel & g_1(c S_1 \cdots c S_n) \xrightarrow{\sim} g_1 c(S_1 + (\cdots + S_n)) & \parallel \text{ product iso for } f \\
 & \downarrow f & \\
 & f S_1 \cdots f S_n &
 \end{array}$$

$$g_1 c S_1 \cdots g_1 c S_n = f S_1 \cdots f S_n$$

Commutates because we are assuming  $f = g_1 c$  as monoidal functors. Now if  $\theta: g_1 \xrightarrow{\sim} g_2$  is such that  $\theta \cdot c: g_1 c \xrightarrow{\sim} g_2 c$  is identity of  $f$ , it follows from the fact that

$$\begin{array}{ccc}
 g_1(c S_1 \cdots c S_n) & \xrightarrow{\sim} & g_1 c(S_1 + \cdots) = f(S_1 + \cdots) \\
 \theta \downarrow & & \theta \downarrow \quad \neq \quad \downarrow \text{id} \\
 g_2(c S_1 \cdots c S_n) & \xrightarrow{\sim} & g_2 c(S_1 + \cdots) = f(S_1 + \cdots)
 \end{array}$$

Thus  $\theta$  is the identity as claimed.



Existence of  $g$ . Define  $g$  on objects:

$$g(c(S_1) \dots c(S_n)) = f_{S_1} \dots f_{S_n}$$

Better we use the ~~map~~ isom.

$$c(S_1) \dots c(S_n) \cong c(S_1 + \dots + S_n)$$

?

Commutativity:

direct product:

If  $\mathcal{S}, \mathcal{T}$  are monoidal categories, then so is  $\mathcal{S} \times \mathcal{T}$  with product functor

$$(\mathcal{S}, \mathcal{T})(\mathcal{S}', \mathcal{T}') = (\mathcal{S}\mathcal{S}', \mathcal{T}\mathcal{T}')$$

~~with~~ with associativity and unity data = the direct product of those of  $\mathcal{S}$  and  $\mathcal{T}$ . Then have strict monoidal functors

$$\mathcal{S} \xleftarrow{\text{pr}_1} \mathcal{S} \times \mathcal{T} \xrightarrow{\text{pr}_2} \mathcal{T}$$

such that

$$\underline{\text{Hom}}^+(\mathcal{U}, \mathcal{S} \times \mathcal{T}) \longrightarrow \underline{\text{Hom}}^+(\mathcal{U}, \mathcal{S}) \times \underline{\text{Hom}}^+(\mathcal{U}, \mathcal{T})$$

is an isomorphism of ~~monoidal~~ categories.

point: The permutal category  $\text{pt}$  has a unique monoidal structure. Claim

$$\underline{\text{Hom}}^+(\text{pt}, \mathcal{S})$$

is the category of pairs  $(\mathcal{S}, 0 \cong \mathcal{S})$ . It is thus a contractible groupoid. In effect a ~~monoidal~~ monoidal functor  $F: \text{pt} \rightarrow \mathcal{S}$  consists of an object  $F(\cdot)$  with product + unity is

$$F(\cdot \cdot) \cong F(\cdot)F(\cdot)$$

$$F(\cdot) \cong 0$$

satisfying conditions of assoc. + unity. The unity condition says:

$$F(\cdot \cdot) \simeq F(\cdot)F(\cdot)$$

$$\parallel \quad \quad \quad |S$$

$$F(\cdot) = 0 F(\cdot)$$

commutes, ~~and~~ which implies  $F(\cdot) \simeq 0$  determines the product isom. On the other hand, it is clear that given  $S \simeq 0$  we ~~can~~ can make  $\square$   $S$  a new 0-object of  $\mathcal{S}$ , hence we get a ~~product isom~~ monoidal functor  $F: \text{pt} \rightarrow \mathcal{S}$  with ~~product isom~~  $F(\cdot) = S$ .

---

There is a distinguished monoidal functor

$$\text{pt} \xrightarrow{F} \mathcal{S}$$

sending  $\cdot$  to  $0$  with  $F(\cdot) \simeq 0 =$  the identity arrow of  $0$ . We will denote this monoidal functor  $O_{\mathcal{S}}: \text{pt} \rightarrow \mathcal{S}$  or simply  $O: \text{pt} \rightarrow \mathcal{S}$ .

---

Problem: To describe  $\text{Hom}^+(\mathcal{S} \times \mathcal{T}, \mathcal{U})$ :

Definition: Let  $F: \mathcal{S} \rightarrow \mathcal{U}$ ,  $G: \mathcal{T} \rightarrow \mathcal{U}$  be monoidal functors. A commutativity-isomorphism between  $F, G$  is an isom

$$F(S)G(T) \simeq G(T)F(S)$$

of functors  $\mathcal{S} \times \mathcal{T} \rightarrow \mathcal{U}$  such that the following conditions hold:

compatible with product iso for F

$$\begin{aligned}
 & F(S_1, S_2) G(T) \simeq G(T) F(S_1, S_2) \\
 & \quad \downarrow \quad \quad \quad \downarrow \\
 & (FS_1, FS_2) G(T) \quad G(T) (FS_1, FS_2) \\
 & \quad \parallel \quad \quad \quad \parallel \\
 & FS_1 (FS_2 G T) \quad (G T \cdot FS_1) FS_2 \\
 & \quad \downarrow \quad \quad \quad \downarrow \\
 & FS_1 (G T / FS_2) = (FS_1 \cdot G T) FS_2
 \end{aligned}$$

compatible with unit iso for F:

$$\begin{aligned}
 & F(0) G(T) \simeq G(T) F(0) \\
 & \quad \downarrow \quad \quad \quad \downarrow \\
 & 0 G(T) \quad G(T) 0 \\
 & \quad \quad \parallel \quad \parallel \\
 & \quad \quad G(T)
 \end{aligned}$$

Similarly ~~the~~ the commutation isom. must be compatible with the product and unit isom. of G.

It is a well-known and useful fact (ascribed to Isbell in [May]) that every monoidal category  $\mathcal{S}$  is equivalent to a strict monoidal category whose objects are elements of the free monoid generated by  $\text{Ob}(\mathcal{S})$ . A complete statement is as follows.

Proposition 2: (i) Given a monoidal category  $\mathcal{S}$  there exists a strict monoidal category and an equivalence of monoidal categories  $\xi: \mathcal{S} \rightarrow \tilde{\mathcal{S}}$  such that  $\text{Ob}(\xi): \text{Ob}(\mathcal{S}) \rightarrow \text{Ob}(\tilde{\mathcal{S}})$  induces an isomorphism of the monoid  $\text{Ob}(\tilde{\mathcal{S}})$  with the free monoid generated by  $\text{Ob}(\mathcal{S})$ . These properties determine the pair  $(\tilde{\mathcal{S}}, \xi)$  up to isom.

(ii) Any monoidal functor  $F: \mathcal{S} \rightarrow \mathcal{T}$ , where  $\mathcal{T}$  is a strict monoidal category, factors uniquely

$$\mathcal{S} \xrightarrow{\xi} \tilde{\mathcal{S}} \xrightarrow{F'} \mathcal{T}$$

where  $F'$  is a strict monoidal functor. Thus the ~~inclusion functor~~ inclusion functor

$$\left\{ \begin{array}{l} \text{strict monoidal categories} \\ \text{strict monoidal functors} \end{array} \right\} \subset \left\{ \begin{array}{l} \text{monoidal cats} \\ \text{monoidal functors} \end{array} \right\}$$

has the left adjoint:  $\mathcal{S} \mapsto \tilde{\mathcal{S}}$ .

One proof of this result is based on MacLane's coherence ~~result~~ <sup>compare</sup> theorem ([May]). We will sketch another proof in § 3.

~~Let  $X$  be a set on which a monoid~~

Let  $S$  be a monoid, and let  $X$  be a set on which  $S$  acts. One can form a category whose objects are the elements of  $X$  and in which a morphism from  $x_1$  to  $x_2$  'is' an element  $s$  of  $S$  such that  $s x_1 = x_2$ . In the following chapter ~~we propose~~ to generalize this construction to the situation ~~of~~

~~we~~ we present an analogue of this construction where  $S$  is replaced by a monoidal category  $\mathcal{S}$  ~~that is a category with a monoidal product and a unit object~~ ~~and~~  $X$  becomes a category  $\mathcal{X}$  on which  $\mathcal{S}$  acts. ~~a monoidal category~~

that is, a category with a product functor which is <sup>(see [Mac Lane])</sup> associative and unitary up to coherent isomorphisms, and where  $X$  becomes a category on which  $\mathcal{S}$  acts.

# monoidal categories and functors

Definition: A monoidal category is a category  $\mathcal{S}$  equipped with

- i) a functor  $\mathcal{S} \times \mathcal{S} \rightarrow \mathcal{S}$  which will be denoted  $(S_1, S_2) \mapsto S_1 S_2$  and called the product of  $\mathcal{S}$
- ii) an object  $0$  of  $\mathcal{S}$  <sup>and</sup> called the unit of  $\mathcal{S}$  which will be denoted  $0_{\mathcal{S}}$  or simply  $0$ ,
- iii) an isomorphism  $(S_1 S_2) S_3 \simeq S_1 (S_2 S_3)$

of functors from  $\mathcal{S} \times \mathcal{S}$  to  $\mathcal{S}$ , called the associativity isom. of  $\mathcal{S}$

$$0 S \simeq S \quad , \quad S 0 \simeq S$$

of functors from  $\mathcal{S}$  to  $\mathcal{S}$ , called the left and right unity isomorphisms, respectively.

The associativity and unity isomorphisms are required to satisfy ~~the~~ <sup>the</sup> ~~compatibility~~ conditions that the diagrams

$$(1) \quad \begin{array}{ccc} S_1(S_2(S_3 S_4)) & \simeq & S_1((S_2 S_3) S_4) \simeq (S_1(S_2 S_3)) S_4 \\ \downarrow S_1 & & \downarrow S_1 \\ (S_1 S_2)(S_3 S_4) & \simeq & ((S_1 S_2) S_3) S_4 \end{array}$$

$$(2) \quad \begin{array}{ccc} (S_1 0) S_2 & \simeq & S_1 (0 S_2) \\ \searrow S_1 & & \swarrow S_1 \\ & S_1 S_2 & \end{array}$$

commute for any ~~the~~ objects  $S_i$  of  $\mathcal{S}$ .

Remarks: From the commutativity of the above diagrams one can deduce the commutativity of

$$\begin{array}{ccc}
 (S_1, S_2) \circ \simeq S_1(S_2 \circ) & & (0S_1)S_2 \simeq 0(S_1, S_2) \\
 \searrow \quad \swarrow & & \searrow \quad \swarrow \\
 & S_1 S_2 & & S_1 S_2
 \end{array}$$

and the fact that the two isomorphisms:  $0 \circ \simeq 0$  obtained from the left and right unity isomorphisms are equal. (see [Kelly]). J. Alg 1 (1964) 377-402

A monoidal category will be called strict if its associativity and unity isomorphisms are identity maps. A strict monoidal category is the same thing as a monoid object in the category of categories, ~~hence the terminology~~ hence ~~the terminology~~ the terminology monoid category will also be used.

If  $\mathcal{X}$  is a category, the category Hom( $\mathcal{X}, \mathcal{X}$ ) of functors from  $\mathcal{X}$  to itself is a strict monoidal category with product given by composition of functors. More generally, the category of endomorphisms of an object in a 2-category is a strict monoidal category.

Definition: Let  $\mathcal{S}, \mathcal{T}$  be monoidal categories. By a monoidal functor  $F: \mathcal{S} \rightarrow \mathcal{T}$  we mean a functor  $S \mapsto FS$  from  $\mathcal{S}$  to  $\mathcal{T}$  equipped with

- i) an isomorphism of functors

$$F(S_1, S_2) \simeq F(S_1)F(S_2)$$

from  $\mathcal{S} \times \mathcal{S}$  to  $\mathcal{T}$ , called the product isom of  $F$

- ii) an isomorphism  $F(0_{\mathcal{S}}) \simeq 0_{\mathcal{T}}$ , called the

unity isomorphism of  $F$ .





Morphism of ~~functors~~ monoidal functors! If

F and G are monoidal functors from S to T, then by a morphism of monoidal functors  $u: F \rightarrow G$ , or monoidal natural transformation, we mean a ~~transformation~~ ~~of the underlying functors~~ morphism of functors which is compatible with the products and unity isomorphisms of F and G. We denote ~~the~~

~~MHom(S, T)~~ ~~monoidal functors~~ ~~from~~ the category <sup>consisting</sup> of monoidal functors ~~from~~ S to T and morphisms between them by MHom(S, T). (MHom is to be script)

Composition of monoidal functors gives rise to functors

$$\underline{\text{MHom}}(\del{S, T}) \times \underline{\text{MHom}}(T, U) \rightarrow \underline{\text{MHom}}(S, U)$$

of the sort that monoidal categories, monoidal functors, and monoidal natural transformations form a 2-category.

Equivalence of monoidal categories: A monoidal

functor  $F: S \rightarrow T$  will be called an equivalence of monoidal categories if there exists a monoidal functor  $G: T \rightarrow S$  and ~~mutual~~ isomorphisms of monoidal functors  $GF \simeq id_S$ ,  $FG \simeq id_T$  which are mutually compatible in the sense that

$$FGF \simeq id_T F$$

$$\downarrow \quad \parallel$$

$$F id_T = F$$

$$GFG \simeq G id_S$$

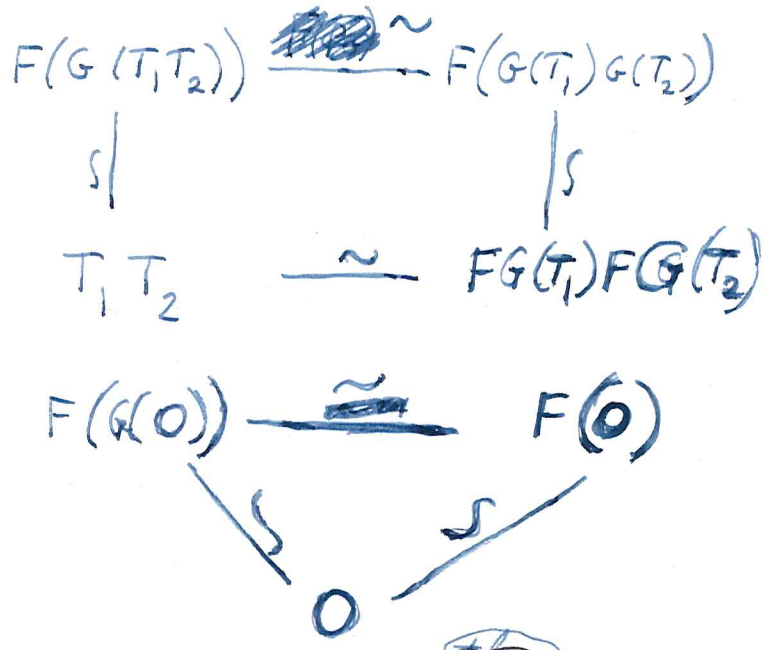
$$\downarrow \quad \parallel$$

$$id_S G = G$$

commute.

Proposition 1: A monoidal functor is an equivalence of monoidal categories if it is an equivalence of the underlying categories.

Proof: Let  $F: \mathcal{S} \rightarrow \mathcal{T}$  be a monoidal functor which is an equivalence of the underlying categories, and let  $G: \mathcal{T} \rightarrow \mathcal{S}$  be a functor quasi-inverse to  $F$  so that there are compatible isomorphisms  $GF \cong id, FG \cong id$ . We equip  $G$  with the unique product isomorphism  ~~$F(G(T_1, T_2)) \cong F(G(T_1)G(T_2))$~~  and ~~unity isom~~ which are compatible with those of  $F$  in the sense that the diagrams



commutes.

~~It is routine to check that  $G$  is a monoidal functor, and that the isos.  $FG \cong id, GF \cong id$  are compatible with the monoidal structures. Therefore  $F$  is an equivalence of monoidal categories.~~

It is routine to check that  $G$  is a monoidal functor, and that the isos.  $FG \cong id, GF \cong id$  are compatible with the monoidal structures. Therefore  $F$  is an equivalence of monoidal categories.

## Commuting monoidal functors.

~~Given monoids  $S, T, U$ , one knows that a monoid homomorphism  $H: S \times T \rightarrow U$  is essentially the same thing~~

~~If  $S, T, U$  are monoids, one knows that to give a monoid homomorphism  $H: S \times T \rightarrow U$  is the same thing as giving a pair of monoid homomorphisms  $F: S \rightarrow U, G: T \rightarrow U$  which commute:  $F(s)G(t) = G(t)F(s)$ .~~

Given monoids  $S, T, U$ , one knows that a monoid homomorphism  $H: S \times T \rightarrow U$  is ~~uniquely~~ expressible in the form  $H(s, t) = F(s)G(t)$  where  $F: S \rightarrow U, G: T \rightarrow U$  are monoid homomorphisms which commute:  $F(s)G(t) = G(t)F(s)$ . We now discuss the generalization of this fact to ~~the~~ monoidal categories.

Definition: Let  $F: S \rightarrow U$  and  $G: T \rightarrow U$  be monoidal functors. By a commutation isomorphism between  $F$  and  $G$  we mean an isom

$$c: F(S)G(T) \xrightarrow{\sim} G(T)F(S)$$

of functors from  $S \times T$  to  $U$  which is compatible with the product + unity isoms. of  $F$  and  $G$ . Compatibility with the product iso. of  $F$  means commutativity of

$$\begin{array}{ccc} F(S_1, S_2)G(T) & \xrightarrow{\sim} & (F(S_1)F(S_2))G(T) \\ \downarrow \cong & & \downarrow \cong \\ G(T)F(S_1, S_2) & & G(T)(F(S_1)F(S_2)) \end{array}$$

$$\begin{array}{ccc}
 F(S_1, S_2) G(T) & \xrightarrow{\sim} & G(T) F(S_1, S_2) \\
 \downarrow s & & \downarrow s \\
 (F(S_1) F(S_2)) G(T) & & G(T) (F(S_1) F(S_2)) \\
 \downarrow s & & \downarrow s \\
 F(S_1) \cancel{G(T)} F(S_2) G(T) & & (G(T) F(S_1)) F(S_2) \\
 \downarrow s & & \downarrow s \\
 F(S_1) (G(T) F(S_2)) & \xrightarrow{\sim} & (F(S_1) G(T)) F(S_2)
 \end{array}$$

while compatibility with the unity iso of  $F$  means that

$$\begin{array}{ccc}
 F(\mathbb{1}) G(T) & \xrightarrow{\sim} & G(T) F(\mathbb{1}) \\
 \downarrow s & & \downarrow s \\
 \mathbb{1} G(T) & \cong & G(T) \cong & G(T) \mathbb{1}
 \end{array}$$

commutes. Compatibility with the product + unity isos of  $G$  is defined similarly.

By abuse of language we say two monoidal functors commute when there is given a commutation isomorphism between them.

Example: We consider the monoidal functor

$$\text{in}_1 = (\text{id}_S, \mathbb{1}) : S \rightarrow S \times T, \quad S \mapsto (S, \mathbb{1})$$

where  $\mathbb{1} : S \rightarrow T$  denotes the monoidal functor given by the constant functor with value  $\mathbb{1}_T$ , equipped with unity iso equal to the identity map of  $\mathbb{1}_T$ , and with product iso furnished by the unity iso  $\mathbb{1}_T \mathbb{1}_T \simeq \mathbb{1}_T$  of  $T$ . Similarly we have the monoidal functor

$$\text{in}_2 = (\mathbb{1}, \text{id}_T) : T \rightarrow S \times T, \quad T \mapsto (\mathbb{1}, T)$$

We claim  
~~that~~, the composite

$$(S, \mathbb{1})(\mathbb{1}, T) = (S\mathbb{1}, \mathbb{1}T)$$

$$\simeq (S, T)$$

$$\simeq (\mathbb{1}S, T\mathbb{1}) = (\mathbb{1}, T)(\mathbb{1}, S)$$

is a commutative isomorphism between  $\text{in}_1$  and  $\text{in}_2$ . To show this we have to ~~show~~<sup>check</sup> the commutativity of certain diagrams constructed from the assoc. & unity isos. of  $S$  and  $T$ . It is clear that we can, without loss of generality, replace  $S$  and  $T$  by equivalent mon. cats. Hence we can suppose  $S, T$  are strict monoidal cats, in which case the relevant diagrams commute because all the arrows occurring in them are identity maps.

~~We are now going to show~~

~~Given a commutation isom  $c$  between  $F: S \rightarrow U$ ,  $G: T \rightarrow U$ , and a monoidal functor  $E: U \rightarrow U'$ , let  $E: U \rightarrow U'$  be a~~

We are now going to show that the above example is in some sense the universal example of a pair of monoidal functors commutation isom between monoidal functors.

Let  $c$  be a commutation isom. between ~~functors~~  $F: S \rightarrow U'$ ,  $G: T \rightarrow U'$ , and let  $H: U' \rightarrow U''$  be a monoidal functor. One verifies easily that the composite isom

~~$$E(F(S))E(G(T)) \simeq E(F(S)G(T)) \xrightarrow{c} E(G(T)F(S)) \simeq E(G(T))E(F(S))$$~~

$$H(F(S))H(G(T)) \simeq H(F(S)G(T))$$

( ) 
$$s \mid H(c)$$

$$H(G(T)F(S)) \simeq H(G(T))H(F(S))$$

is ~~is~~ a commutation isomorphism between  $EF: S \rightarrow U''$  and  $EG: T \rightarrow U''$ . In particular taking  $U' = S \times T$  and  $F, G$  to be ~~the functions~~  $in_1, in_2$ , respectively, we have associated to any monoidal functor  $H: S \times T \rightarrow U$  a pair of monoidal functors

$$\begin{array}{l}
 H_{in_1}: S \rightarrow U \\
 H_{in_2}: T \rightarrow U
 \end{array}
 , \quad
 \begin{array}{l}
 S \mapsto H(S, \mathbb{1}) \\
 T \mapsto H(\mathbb{1}, T)
 \end{array}$$

equipped with the commutation isomorphism:

$$\begin{aligned} H(S, \mathbb{1}) H(\mathbb{1}, T) &\simeq H(S\mathbb{1}, \mathbb{1}T) \\ &\simeq H(S, T) \\ &\simeq H(\mathbb{1}S, T\mathbb{1}) \\ &\simeq H(\mathbb{1}, T) H(S, \mathbb{1}). \end{aligned}$$

~~We wish to show that ~~any~~ up to isom. a monoidal functor ~~H~~  $H: S \times T \rightarrow \mathcal{U}$  may be identified with a pair of mon. functors  $F: S \rightarrow \mathcal{U}$ ,  $G: T \rightarrow \mathcal{U}$  and a commutation isom between them. To formulate this precisely we ~~introduce~~ introduce the category  $\mathcal{H}(S, T; \mathcal{U})$  defined as follows. An object is a triple  $(F, G, c)$  where  $F$~~

We ~~wish~~ wish to show that in this way a mon. functor  $H: S \times T \rightarrow \mathcal{U}$  may be identified, up to isomorphism, with a triple  $(F, G, c)$  consisting of monoidal functors  $F: S \rightarrow \mathcal{U}$ ,  $G: T \rightarrow \mathcal{U}$  and a comm. isom  $c$  between them. To formulate this precisely, let  $\mathcal{H}(S, T; \mathcal{U})$  denote the category of such triples in which a morphism  $(F, G, c) \rightarrow (F', G', c')$  consists of two morphisms  $F \rightarrow F'$ ,  $G \rightarrow G'$  of monoidal functors which are compatible with  $c$  and  $c'$  in the evident sense. Then it is clear that by associating to  $H: S \times T \rightarrow \mathcal{U}$  the pair  $H_{in_1}, H_{in_2}$  with the commutation iso described above, we obtain a functor

$$(\ ) \quad \underline{M\text{Hom}}(S \times T, \mathcal{U}) \longrightarrow \mathcal{H}(S; T; \mathcal{U}).$$

The desired result may now be formulated as follows:



Proposition: The functor  $( )$  is an equivalence of categories.

To prove this we construct a functor in the opposite direction. Let  $(F, G, c)$  be an object of  $\mathcal{H}(S, T; \mathcal{U})$  and let  $H$  be the functor from  $S \times T$  to  $\mathcal{U}$  given by

$$H(S, T) = F(S) G(T).$$

Define a product isom for  $H$  to be the composition

$$\begin{aligned} H((S_1, T_1)(S_2, T_2)) &= H(S_1 S_2, T_1 T_2) \\ &= F(S_1 S_2) G(T_1 T_2) \\ &\cong (F(S_1) F(S_2)) (G(T_1) G(T_2)) \\ &\cong F(S_1) [F(S_2) (G(T_1) G(T_2))] \\ &\cong F(S_1) [(F(S_2) G(T_1)) G(T_2)] \\ &\cong F(S_1) [(G(T_1) F(S_2)) G(T_2)] \\ &\cong F(S_1) [G(T_1) (F(S_2) G(T_2))] \\ &\cong (F(S_1) G(T_1)) (F(S_2) G(T_2)) \\ &= H(S_1, T_1) H(S_2, T_2). \end{aligned}$$

Define the unity isomorphism for  $H$  to be the composite

$$H(1, 1) = F(1) G(1) \simeq 1 \cdot 1 \simeq 1$$

Lemma: The functor  $H$ , <sup>equipped</sup> with the above product and unity isomorphisms is a monoidal functor from  $S \times T$  to  $\mathcal{U}$ .

Proof. We first reduce to the case where  $S, T, \mathcal{U}$  are strict monoidal categories and  $F, G$  are strict monoidal functors. Notice that to prove the lemma we have only to show

the commutativity of certain diagrams in  $\mathcal{U}$  constructed from the ~~monoidal~~ <sup>comm. isom and</sup> ~~isomorphisms~~ belonging to the monoidal structure of  $S, T, \mathcal{U}, F, G$ . It is clear that if two triples  $(F, G, c)$  and  $(F', G', c')$  are isomorphic, then the lemma holds for one iff it does for the other. Further if the lemma holds for  $(F, G, c)$  then it also does for the triples obtained by pulling back with respect to monoidal functors  $S' \rightarrow S, T' \rightarrow T$ , and pushing forward with respect to  $\mathcal{U} \rightarrow \mathcal{U}'$ . Thus it is clear that ~~we can replace~~ ~~without loss of generality~~ we can replace  $S, T, \mathcal{U}, F, G$  by  $\tilde{S}, \tilde{T}, \tilde{\mathcal{U}}, \tilde{F}, \tilde{G}$  and so reduce to the strict cases.

Lemma 1: The functor  $H$  equipped with the above prod + unly isos is a monoidal functor from  $S \times T$  to  $\mathcal{U}$ .

Proof: ~~To prove this lemma we must show that certain~~ the commutativity of certain diagrams in  $\mathcal{U}$  constructed from the comm. iso  $c$  and the isos. given by the monoidal structure of  $S, T, \mathcal{U}, F, G$ . Since the canonical functor  $\gamma: \mathcal{U} \rightarrow \tilde{\mathcal{U}}$  is an equivalence of monoidal categories, it ~~evidently~~ evidently suffices to prove the lemma for the ~~triple~~ ~~triple~~  $\gamma F: S \rightarrow \tilde{\mathcal{U}}, \gamma G: T \rightarrow \tilde{\mathcal{U}}$  and the ~~triple~~ triple consisting of  $\gamma F: S \rightarrow \tilde{\mathcal{U}}, \gamma G: T \rightarrow \tilde{\mathcal{U}}$  and the commutation isom <sup>between them</sup> induced from  $c$ . Since  $S \rightarrow \tilde{S}, T \rightarrow \tilde{T}$  are equivalences of mon. cats, one sees easily that this induced comm. iso ~~extends to induce~~ <sup>extends to induce</sup> a commutation isom  $\tilde{c}$  between  $\tilde{F}: \tilde{S} \rightarrow \tilde{\mathcal{U}}, \tilde{G}: \tilde{T} \rightarrow \tilde{\mathcal{U}}$ . It therefore suffices to ~~prove~~ prove the lemma for  $(\tilde{F}, \tilde{G}, \tilde{c})$ . Thus without loss of generality, we can suppose  $S, T, \mathcal{U}, F, G$  are strict monoidal cats + functors.

In this case the product isom for  $H(S, T) = F(S) G(T)$  becomes

$$\begin{aligned}
 H((S_1, T_1)(S_2, T_2)) &= FS_1 FS_2 GT_1 GT_2 \\
 &\cong FS_1 GT_1 FS_2 GT_2 = H(S_1, T_1) H(S_2, T_2)
 \end{aligned}$$

and the condition expressing compatibility of the product isomorphism with associativity isomorphisms ~~is~~ is for

$$H((S_1, T_1)(S_2, T_2)(S_3, T_3)) \sim H((S_1, T_1)(S_2, T_2)) H(S_3, T_3)$$

$$\quad \quad \quad \Big| \quad \quad \quad \quad \quad \quad \quad \Big|$$

$$H(S_1, T_1) H((S_2, T_2)(S_3, T_3)) \sim H(S_1, T_1) H(S_2, T_2) H(S_3, T_3)$$

to commute. However this is easily seen to be the outside square of the diagram in Figure 1.

$$\begin{array}{ccc}
 F(S_1 S_2 S_3) G(T_1 T_2 T_3) = F(S_1 S_2) F S_3 G(T_1 T_2) G T_3 & \sim & F(S_1 S_2) G(T_1 T_2) F S_3 G T_3 \\
 \parallel & & \parallel \\
 F S_1 F(S_2 S_3) G T_1 G(T_2 T_3) = F S_1 F S_2 F S_3 G T_1 G T_2 G T_3 & & \\
 \Big| & & \Big| \\
 & & F S_1 F S_2 G T_1 F S_3 G T_2 G T_3 \sim F S_1 F S_2 G T_1 G T_2 F S_3 G T_3 \\
 & & \Big| \quad \quad \quad \Big| \\
 F S_1 G T_1 F(S_2 S_3) G(T_2 T_3) & \sim & F S_1 G T_1 F S_2 F S_3 G T_2 G T_3 \sim F S_1 G T_1 F S_2 G T_2 F S_3 G T_3
 \end{array}$$

Figure 1.

~~The axioms on a commutative isomorphism~~ The axioms on a commutative isomorphism imply the pentagons in this diagram commute. ~~The~~ The lower right square commutes because the product is a bifunctor of its arguments. Hence the diagram commutes as desired.

~~In the strict case the unit isomorphism for H is the identity map of U, and by the axiom that the comm. isom. is~~

In the case where  $F, T, U, F, G$  are strict, the condition that  $c$  be compatible with the unit isomorphism of  $F$  ~~is~~ asserts that the comm. isom.

$$F(1)G(T) \simeq G(T)F(1)$$

is the identity map of  $G(T)$ . The unit isom for  $H$  is the identity map of  $H(1,1) = 1$ . Hence all the arrows in

$$H((1,1)(S,T)) \simeq H(1,1)H(S,T)$$

$\downarrow$

$\downarrow$

$$H(S,T) \simeq 1H(S,T)$$

are equal to the identity map of  $F(S)G(T)$ . Thus ~~the~~  <sup>$H$  satisfies</sup> the left unity condition for a monoidal functor, ~~and the~~. The right one is similar, so the lemma is proved.

By virtue of the lemma, it is clear that by associating to  $(F,G,c)$  the <sup>monoidal</sup> functor  $H$ , we obtain a functor

$$(**) \quad H(S,T; \mathcal{U}) \longrightarrow \underline{MHan}(S \times T; \mathcal{U}).$$

Lemma 2: The functors  $(**)$  and  $(*)$  are quasi-inverses of each other.

~~Proof.~~

~~Proof.~~

~~Under the composition  $(**)(*)$  a mon. functor~~

~~$H: S \times T \rightarrow \mathcal{U}$  goes into the functor~~

$$H'(S,T) = H(S, \mathbb{1})H(\mathbb{1}, T)$$

~~equipped with certain product and unity isos.~~

The reader may <sup>verify</sup> ~~the~~

~~isomorphism of functors~~

$$H(S, \mathbb{1})H(\mathbb{1}, T) \simeq H(S, \mathbb{1}T) \simeq H(S, T)$$

~~is an isomorphism of monoidal functors,~~

~~compatible with product and unity isos, and further that~~ <sup>in this way</sup> one obtains an isomorphism of the composite  $(**)(*)$  with the identity. 10

Lemma 2: The functors  $(**)$  and  $(*)$  are quasi-inverses of each other.

Proof. Under the composite functor  $(**)(*)$  a monoidal functor  $H: S \times T \rightarrow \mathcal{U}$  goes into the functor

$$(S, T) \mapsto H(S, 1)H(1, T),$$

equipped with certain product + unity isomorphisms. The reader may verify that the isomorphism

$$H(S, 1)H(1, T) \simeq H(S1, 1T) \simeq H(S, T)$$

is an isomorphism of monoidal functors (by considerations of naturality one reduces to the case where  $H$  is the identity functor of  $S \times T$ , in which case the verification is almost trivial). It is clear that in this way the ~~composite~~ <sup>functor</sup>  $(**)(*)$  is isom. to the identity.

Under the composite functor  $(*)(**)$  a triple  $(F, G, c)$  goes to the pair of functors

$$S \mapsto F(S)G(1)$$

$$T \mapsto F(1)G(T)$$

equipped with certain monoidal structures and a commutation isomorphism. The reader may verify that the isos.

$$F(S)G(1) \simeq F(S)1 \simeq F(S)$$

$$F(1)G(T) \simeq 1G(T) \simeq G(T)$$

constitute isomorphisms of monoidal functors compatible with the commutations isos. (~~Verification~~ One ~~reduces~~ reduces to the case where  $S, T, \mathcal{U}, F, G$  are strict, in which case the verification is trivial.) It is thus clear that  $(*)(**)$  is isomorphic to the identity, and so the lemma follows.

# Commuting actions

~~Let  $S, T$  be monoidal ~~functors~~ <sup>categories</sup> and let  $\alpha$   
 $\alpha: S \rightarrow \underline{\text{Hom}}(\mathcal{X}, \mathcal{X})$   
 $\beta: T \rightarrow \underline{\text{Hom}}(\mathcal{X}, \mathcal{X})$   
 be actions of  $S$  and  $T$  on a category  $\mathcal{X}$~~

Let two monoidal categories  $S, T$  act on the same category  $\mathcal{X}$ , ~~the~~ the actions being defined by mon. functors

$$\alpha: S \rightarrow \underline{\text{Hom}}(\mathcal{X}, \mathcal{X}), \quad S \mapsto (X \mapsto SX)$$

$$\beta: T \rightarrow \underline{\text{Hom}}(\mathcal{X}, \mathcal{X}), \quad T \mapsto (X \mapsto TX)$$

A commutation isom. between  $\alpha$  and  $\beta$  may be identified with an isomorphism

$$S(TX) \simeq T(SX)$$

of functors from  $S \times T \times \mathcal{X}$  to  $\mathcal{X}$ , which is compatible <sup>in the evident sense</sup> with the associativity and unity isos. for the  $S$  and  $T$  action. It may also be identified with a lifting of  $\alpha$  to a monoidal functor

$$\alpha': S \rightarrow \underline{\text{Hom}}^T(\mathcal{X}, \mathcal{X})$$

or with a lifting of  $\beta$  to a monoidal functor

$$\beta': T \rightarrow \underline{\text{Hom}}^S(\mathcal{X}, \mathcal{X})$$

By abuse of language, we will say that the  $S$  and  $T$  actions ~~commute~~ commute if there is given a commutation isomorphism between  $\alpha$  and  $\beta$ .

On composing  $\alpha'$  with the canonical functor

$$\underline{\text{Hom}}^T(\mathcal{X}, \mathcal{X}) \rightarrow \underline{\text{Hom}}(\langle T, \mathcal{X} \rangle, \langle T, \mathcal{X} \rangle) \quad \text{AOD:}$$

$\underline{\text{Hom}}^S(\mathcal{X}, \mathcal{X}) \rightarrow \underline{\text{Hom}}(\langle S, \mathcal{X} \rangle, \langle S, \mathcal{X} \rangle)$

we obtain an <sup>induced</sup> action of  $S$  on  $\langle T, X \rangle$ . and on  $T^{-1}X$   
 there is an <sup>induced</sup> action of  $T$  on  $\langle S, X \rangle$ . and on  $S^{-1}X$ . Similarly

~~hand from the preceding section (see lemma 1) there is a monoidal functor~~

~~$$S \times T \longrightarrow \underline{\text{Hom}}(X, X), \quad (S, T) \longmapsto (X \rightarrow S(TX))$$~~

~~with product isomorphisms~~

~~$$(S_1, T_1)$$~~

On the other hand the commutation isom between  $\alpha$  and  $\beta$  gives rise to a monoidal functor

$$S \times T \longrightarrow \underline{\text{Hom}}(X, X)$$

(see lemma 1 ...), hence to an action of  $S \times T$  on  $X$ . This action is given by

$$(S, T)X = S(TX)$$

with the associativity isom

$$\begin{aligned} ((S_1, S_2)(T_1, T_2))X &= (S_1, S_2)((T_1, T_2)X) \\ &\simeq S_1(S_2(T_1(T_2X))) \\ &\simeq \del{S_1(T_1(S_2(T_2X)))} S_1(T_1(S_2(T_2X))) \\ &= (S_1, T_1)((S_2, T_2)X) \end{aligned}$$

and unity isom

$$(0, 0)X = 0(0X) \simeq 0X \simeq X.$$

Proposition: (i) The categories  $\langle S, \langle T, X \rangle \rangle$ ,  $\langle T, \langle S, X \rangle \rangle$  and  $\langle S \times T, X \rangle$  are canonically isom.

(ii) The categories  $S^{-1}(T^{-1}X)$ ,  $T^{-1}(S^{-1}X)$ , and  $(S \times T)^{-1}X$  are canonically isomorphic.

Proof. (i): The categories in question have the same objects as  $\mathcal{X}$ . ~~On the other hand~~ As

~~the following~~

$$\begin{aligned} \text{Hom}_{\langle \mathcal{S} \times \mathcal{T}, \mathcal{X} \rangle}(X, X') &= \lim_{(\mathcal{S}, \mathcal{T})} \text{Hom}_{\mathcal{X}}((\mathcal{S}, \mathcal{T})X, X') \\ &\cong \lim_{\mathcal{T}} \lim_{\mathcal{S}} \text{Hom}_{\mathcal{X}}(\mathcal{S}(\mathcal{T}X), X') \\ &= \lim_{\mathcal{T}} \text{Hom}_{\langle \mathcal{S}, \mathcal{X} \rangle}(\mathcal{T}X, X') \\ &= \text{Hom}_{\langle \mathcal{T}, \langle \mathcal{S}, \mathcal{X} \rangle \rangle}(X, X') \end{aligned}$$

one obtains ~~there is~~ a 1-1 correspondence between arrows in the cats  $\langle \mathcal{S} \times \mathcal{T}, \mathcal{X} \rangle$  and  $\langle \mathcal{T}, \langle \mathcal{S}, \mathcal{X} \rangle \rangle$  ~~which~~ <sup>by associating to</sup> ~~the~~  $\langle \mathcal{S} \times \mathcal{T}, \mathcal{X} \rangle$  ~~map~~ <sup>-map</sup> represented by ~~the~~  $\mathcal{X}$ -map  $x: (\mathcal{S}, \mathcal{T})X \rightarrow X'$

one obtains a 1-1 correspondence between arrows in the cats  $\langle \mathcal{S} \times \mathcal{T}, \mathcal{X} \rangle$  and  $\langle \mathcal{T}, \langle \mathcal{S}, \mathcal{X} \rangle \rangle$  <sup>which may be described as follows</sup>: ~~the~~  $\langle \mathcal{S} \times \mathcal{T}, \mathcal{X} \rangle$ -map ~~the~~  $X \rightarrow X'$  represented by the  $\mathcal{X}$ -map  $x: (\mathcal{S}, \mathcal{T})X \rightarrow X'$  corresponds the  $\langle \mathcal{T}, \langle \mathcal{S}, \mathcal{X} \rangle \rangle$ -map represented by the  $\langle \mathcal{S}, \mathcal{X} \rangle$ -map  $\mathcal{T}X \rightarrow X'$  rep. by the  $\mathcal{X}$ -map  $x: \mathcal{S}(\mathcal{T}X) \rightarrow X'$ . Using this description, it is easy to see that the 1-1 correspondence is compatible with composition. Thus the cats  $\langle \mathcal{S} \times \mathcal{T}, \mathcal{X} \rangle$  and  $\langle \mathcal{T}, \langle \mathcal{S}, \mathcal{X} \rangle \rangle$  are isomorphic. ~~and by symmetry~~  $\langle \mathcal{S}, \langle \mathcal{T}, \mathcal{X} \rangle \rangle$  is isomorphic to  $\langle \mathcal{S} \times \mathcal{T}, \mathcal{X} \rangle$ .

(ii). Using (i) and (ref. —) we have

$$(\mathcal{S} \times \mathcal{T})^{-1} \mathcal{X} = \langle \mathcal{S} \times \mathcal{T}, \mathcal{S} \times \mathcal{T} \times \mathcal{X} \rangle$$

$$\cong \langle \mathcal{S}, \langle \mathcal{T}, \mathcal{S} \times \mathcal{T} \times \mathcal{X} \rangle \rangle$$

$$\cong \langle \mathcal{S}, \mathcal{S} \times \langle \mathcal{T}, \mathcal{T} \times \mathcal{X} \rangle \rangle = \mathcal{S}^{-1} \mathcal{T}^{-1} \mathcal{X}.$$

Similarly for  $\mathcal{T}^{-1}(\mathcal{S}^{-1} \mathcal{X})$ .

QED



At the moment the writing process is stuck on ~~some aspects~~ of the  $\mathcal{S}$ - $\mathcal{P}$ -section

monoidal cat = cat  $\mathcal{S}$  equipped with product functor, associativity isom, unit object, and left & right unit isos.  $\exists$  various compatibilities hold

example: strict monoidal cats

monoidal functor = functor equipped with a product isom + a unit isom subject to compat with the assoc. + left & right-unit isos.

strict monoidal functors  
composition of monoidal functors

~~morphism of monoidal functors~~  
morphism of monoidal functors = morphism of underlying functors comp. with prod. + unity iso.

$\text{Hom}^*(\mathcal{S}, \mathcal{T})$  and 2-cat of monoidal

~~the~~ equivalence of monoidal categories + Lemma

(left) action of a monoidal cat.  $\mathcal{S}$  on a category  $\mathcal{X}$

= monoidal functor  $\mathcal{S} \rightarrow \text{Hom}(\mathcal{X}, \mathcal{X})$

=  $\left\{ \begin{array}{l} \text{product} \\ \text{assoc iso} \\ \text{unit iso} \end{array} \right. \begin{array}{l} (\mathcal{S}, \mathcal{X}) \mapsto \mathcal{S}\mathcal{X} \\ (\mathcal{S}_1, \mathcal{S}_2)\mathcal{X} \simeq \mathcal{S}_1(\mathcal{S}_2\mathcal{X}) \\ 0\mathcal{X} \simeq \mathcal{X} \end{array} +$

Example: ~~left~~ actions of  $\mathcal{S}, \mathcal{S}^*$  on  $\mathcal{S}$ .

action-preserving functor  $F: \mathcal{X} \rightarrow \mathcal{Y} =$

functor equipped with an action-preserving isom

$$F(\mathcal{S}\mathcal{X}) \simeq \mathcal{S}F(\mathcal{X})$$

comp. with assoc. + unity isos.

composition

morphism of action-preserving functors

2-cat of cats ~~with~~ on which  $\mathcal{S}$  ~~open~~ acts.

$F(\mathcal{S}_1, \mathcal{S}_2)$