

Let N be a countable infinite set, and G the group of all permutations of N . I want to show G is acyclic, i.e. $\tilde{H}_*(G, \mathbb{Z}) = 0$.

First part: The "infinite sum" argument.

Let N be the groupoid consisting of all countable infinite sets and bijections between them. One has $H_*(G) = H_*(N)$. The disjoint union functor

$$N \times N \xrightarrow{\sqcup} N$$

induces a product $\mu: H_*(G) \otimes H_*(G) \rightarrow H_*(G)$

~~which is associative and commutative. If e is the obvious generator of $H_0(G) = \mathbb{Z}$, then multiplication by $e: H_*(G) \rightarrow H_*(G)$ $x \mapsto e \cdot x$ is the map induced by the embedding $G \hookrightarrow G$ which results from an embedding $N \hookrightarrow N$ with infinite complement.~~

Let $\Sigma: N \rightarrow N$, $X \mapsto N \times X$ denote the infinite sum functor. One has

$$\Sigma \dashv \text{id} \simeq \Sigma,$$

i.e. the diagram

$$\begin{array}{ccccc} n & \xrightarrow{\Delta} & n \times n & \xrightarrow{\Sigma \text{id}} & n \times n \xrightarrow{\pi} n \\ & & \underbrace{\hspace{10em}}_{\Sigma} & & \end{array}$$

commutes up to isomorphism. Let $\alpha \in \tilde{H}_*(G)$ and

$$(\Delta_*)(\alpha) = \alpha \otimes e + \sum_i \alpha'_i \otimes \alpha''_i + e \otimes \alpha$$

with $\deg(\alpha'_i), \deg(\alpha''_i) < \deg(\alpha)$. (I have to pass to mod p or rational homology to have such a formula.) Then since

$$\Sigma_* = \mu(\Sigma_* \otimes \text{id}) \Delta_*$$

one has

$$\Sigma_* \alpha = \Sigma_* \alpha \cdot e + \sum_i \Sigma_* \alpha'_i \cdot \alpha''_i + \Sigma_* e \cdot \alpha$$

Assuming that $e \cdot \beta = 0$ for $0 < \deg(\beta) < \deg(\alpha)$, one gets

$$e \cdot \Sigma_* \alpha = \Sigma_* \alpha \cdot e^2 + 0 + e \cdot \Sigma_* e \cdot \alpha.$$

As $e^2 = e$, $\Sigma_* e = e$, one gets $e \cdot \alpha = 0$. Thus:

Lemma

1: Multiplication by e on $\tilde{H}_*(G)$ is zero.

Second part: The "building" argument:

Let X be the following poset. An element

of X is a subset $x \subset N$ such that both x and $N-x$ are infinite. One has $x \leq y$ in X iff either $x=y$ or $x \subset y$ and $y-x$ is infinite.

Lemma 2: X is contractible.

Proof: Let F be a finite subset of X .

I am going to produce an element z' of X such that ~~such that $x \leq z'$ for all $x \in F$~~

a) $x \mapsto x \vee z'$ is a morphism of posets from F to X , b) $x \leq x \vee z \geq z'$ for all $x \in F$. It follows that the inclusion functor $F \hookrightarrow X$ is homotopic to the constant functor ~~with value~~ with value z' . Thus F contracts to a point in X , so X is contractible (every element of $\pi_*(X)$ comes from some finite F).

Choose a maximal subset x_1, \dots, x_l of F such that $y = x_1 \cup \dots \cup x_l$ is infinite. For any x in F not equal to one of the x_i , we have $x \vee y$ is finite, so ~~is~~ after removing each of these finite subsets from y we obtain an element z of X such that for any x in F , either $z \subset x$ or $z \cap x = \emptyset$.

Divide z into two infinite pieces $z = z' \sqcup z''$. Let's check z' satisfies a) and b). Let $x \in F$. If

$z < x$, then $x \cup z' = x \in F$ and $x \leq x \cup z'$.
 Also $z' \leq x \cup z'$ because $z'' < (x \cup z') - z'$.
 On the other hand $z \cap x = \emptyset$, then $x \cup z' \in F$ because
 its complement contains z'' . Also $x \leq x \cup z' \geq z'$ in
 this case. This proves b) and part of a).

It remains to show that if $x_1 \leq x_2$ in F , then
 $x_1 \cup z' \leq x_2 \cup z'$ in F . This is clear if either $z < x_1$,
 or if $z \cap x_2 = \emptyset$. Suppose then that we take the
 remaining case $x_1 \cap z = \emptyset$, $\blacksquare z < x_2$. Then $x_1 \cup z' \leq x_2$
 $= x_2 \cup z'$ because of z'' . Q.E.D.

Now we consider the complex of chains on X :

$$\rightarrow \bigoplus_{x_0 \leq \dots \leq x_p} \mathbb{Z} \rightarrow \dots \rightarrow \bigoplus_{x_0 \leq x_1} \mathbb{Z} \rightarrow \bigoplus_{x_0} \mathbb{Z} \rightarrow \mathbb{Z} \rightarrow 0$$

which is acyclic because X is contractible. The
 group G acts on X , ~~hence~~ hence the above
 sequence is a complex of G -modules. It gives rise
 to a spectral sequence:

$$E^1_{pq} = H_q(G, \bigoplus_{x_0 \leq \dots \leq x_{p-1}} \mathbb{Z}) \implies 0$$

Note that if $x_0 \leq \dots \leq x_p$ is a p -simplex $\overset{\text{in}}{\in} X$, then one

has ~~the~~ a decomposition $N = N_0 \sqcup \dots \sqcup N_{p+1}$ ~~the~~
 with $x_j = N_0 \sqcup \dots \sqcup N_j$ and where each N_j is infinite.
 It follows that G acts transitively on the set of
 p -simplices and that the stabilizer of a typical
 simplex is isomorphic to $G \times \dots \times G$ ($p+2$ times). \blacksquare
~~Therefore~~ Thus the group of $(p-1)$ chains
 is an induced module:

$$\bigoplus_{x_0 < \dots < x_{p-1}} \mathbb{Z} \hookrightarrow \mathbb{Z}[G] \otimes_{\mathbb{Z}[G^{p+1}]} \mathbb{Z}$$

and so by Shapiro's lemma:

$$H_*(G, \bigoplus_{x_0 < \dots < x_{p-1}} \mathbb{Z}) \simeq H_*(G^{p+1}, \mathbb{Z}).$$

~~This is good~~

It appears that I should be working with
 coefficients ~~ring a field~~ F , whence by Künneth:

$$(*) \quad H_*(G, \bigoplus_{x_0 < \dots < x_p} F) \simeq H_*(G^p, F) \simeq H_*(G, F)^{\otimes (p+1)}$$

The differential d_i in the spectral sequence is the
 alternating sum of the "face" maps

$$d_i : H_*(G, \bigoplus_{x_0 < \dots < x_p} F) \longrightarrow H_*(G, \bigoplus_{x_0 < \dots < x_{p-2}} F)$$

which one obtains by deleting x_i from the simplex

$x_0 < \dots < x_p$. Under the Shapiro isomorphism d_i will correspond to the map $G^{p+1} \rightarrow G^p$ sending $(g_1, \dots, g_i, g_{i+1}, \dots, g_p) \mapsto (g_1, \dots, g_i \oplus g_{i+1}, \dots, g_p)$, where $\oplus: G \times G \rightarrow G$ is the map obtained from any isom. $N \oplus N \cong N$.

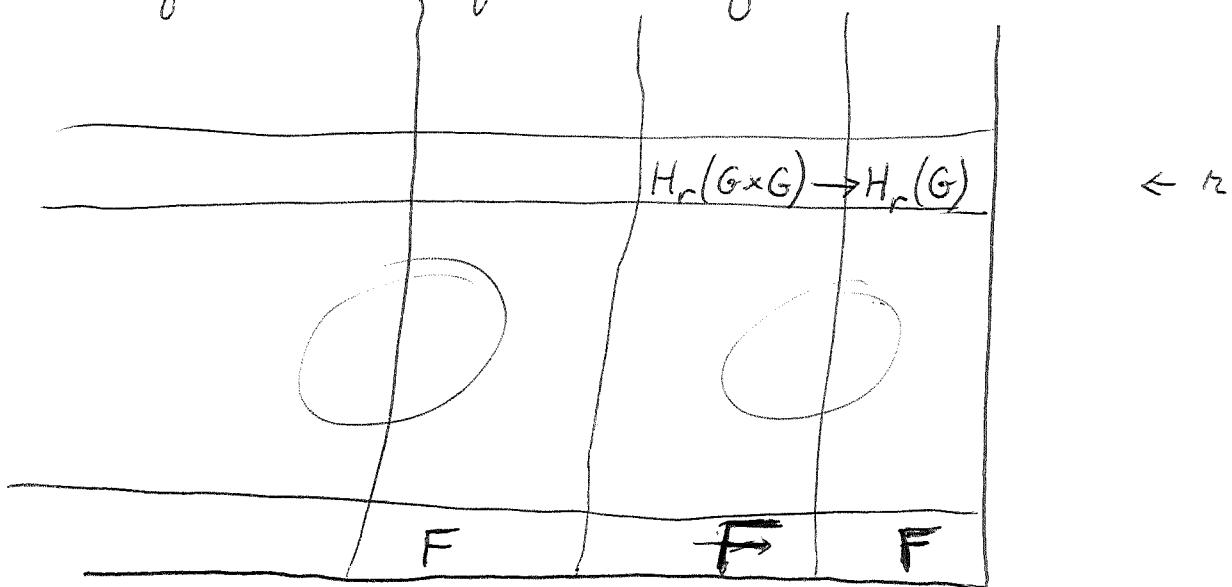
One sees that $E_{0p}^1 = H_0(G^{p+1}, F) \cong F$

with $d_i = \text{id}$, hence $d = d_0 - d_1 + \dots \pm d_p = \text{id}$ if p is odd and 0 if p is even. So E_{0*}^1 is

$$\dots \rightarrow F \xrightarrow{1} F \xrightarrow{0} F \xrightarrow{1} F$$

and $E_{0*}^2 = 0$.

I want to show that $\tilde{H}_*(G, F) = 0$. Suppose that $\tilde{H}_g(G, F) = 0$ for all $g < n$.



By the spectral sequence the map $H_r(G \times G) \xrightarrow{d} H_r(G)$ must be onto. But $H_r(G \times G) = H_r(G) \oplus H_r(G)$, so one

sees that the map

$$\begin{aligned} H_r(G) \oplus H_r(G) &\longrightarrow H_r(G) \\ \alpha &\quad \beta \longmapsto \alpha \cdot e + e \cdot \beta \end{aligned}$$

is onto. But by Lemma 1, this map is 0, so $H_r(G) = 0$

so ~~_____~~ I have proved that $\tilde{H}_*(G, F) = 0$
for all fields F . This implies $\tilde{H}_*(G, \mathbb{Z}) = 0$.

Problem: Let K be a group. Then G acts
on $K^N = \{\text{functions: } N \rightarrow K\}$, so we can form the
group $G \times K^N$ = semi-direct product. (This is
a kind of infinite wreath product $\ast G \wr K$). Is
the group $G \times K^N$ acyclic? Note: Kan + Thurston
have shown this when instead of K^N one takes maps
 $\mathbb{Q} \rightarrow N$ with compact support, and instead of G
one takes maps $\mathbb{Q} \rightarrow \mathbb{Q}$ with compact support.

of the group in the sense of homological algebra

Homology of BC. It is well-known that the homology of the classifying space of a discrete group ~~coincides with the~~ coincides with the ~~homology~~ ~~straightforward~~ homology. We ~~can~~ now describe the generalization of this fact to an arbitrary small category which is perhaps less well-known.

$\delta \xrightarrow{\lambda} \bar{\delta} = \underline{\text{Horn}}^{\delta^*}(\delta, \delta)$ is an equiv. of monoidal categories

$$\begin{array}{ccc} \delta & \xrightarrow{F} & \mathcal{U} \\ \downarrow \delta & \nearrow \delta' & \downarrow H \\ \bar{\delta} & \xrightarrow{G} & \bar{\mathcal{U}} \end{array}$$

$$\delta': (\mathcal{G}, \mathcal{U}, FG(0) \simeq \mathcal{U})$$

$$\delta \xrightarrow{\lambda} \delta' \quad S \mapsto (\lambda_S, F(S), F(\lambda_S(0)) \simeq F(S)).$$

organize as follows.

Statement: Given a monoidal cat \mathcal{S} there exists a strict mon. categ. $\tilde{\mathcal{S}}$ and an equiv of monoidal

$\tilde{\mathcal{S}}$ characterized up to isom by fact that it is strict & that $\exists \tilde{F}: \mathcal{S} \rightarrow \tilde{\mathcal{S}}$ mon. equiv. $\Rightarrow \text{Ob}(\tilde{\mathcal{S}}) = \text{free monoid gen. by } \text{Ob}(\mathcal{S})$.

$\mathcal{S} \xrightarrow{\sim} \tilde{\mathcal{S}}$ left adjoint. Any $F: \mathcal{S} \rightarrow \mathcal{U}$ with \mathcal{U} strict factors uniquely $F = F' \circ \tilde{F}$, where $F': \tilde{\mathcal{S}} \rightarrow \mathcal{U}$ is a strict mon. functor

~~skipped by~~ Constructing a mon. equiv $d: \mathcal{S} \rightarrow \tilde{\mathcal{S}}$ with \mathcal{S} strict.

Construction of $\mathcal{D} \xrightarrow{F} \mathcal{U}$

Proof: Suppose given $d: \mathcal{S} \rightarrow \tilde{\mathcal{S}}$ with \mathcal{S} strict. If S is a set - let $w(S)$ be the cat ~~whose objects are~~ of sort that a monoid goes to a monoid cat. Then put

$$\tilde{\mathcal{S}} = w(F.M.(\text{Ob}(\mathcal{S}))) \times \tilde{\mathcal{S}}$$

In other words $\text{Ob}(\tilde{\mathcal{S}})$

Note that if $F.M(\text{Ob}(\mathcal{S})) \rightarrow \text{Ob}(\tilde{\mathcal{S}})$ is an isom $\Leftrightarrow \tilde{\mathcal{S}} \rightarrow \tilde{\mathcal{S}}$

Lemma 2: Let \mathcal{T}, \mathcal{U} be strict monoidal categories and $F: \mathcal{T} \rightarrow \mathcal{U}$ a mon. functor.

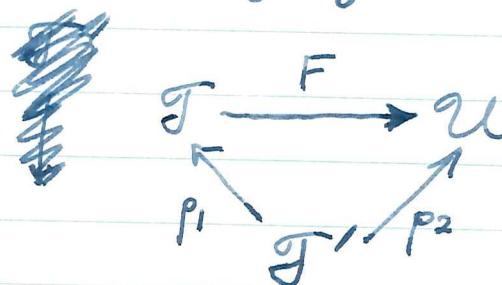
Construction: Let \mathcal{T}, \mathcal{U} be strict monoidal categories and let $F: \mathcal{T} \rightarrow \mathcal{U}$ be a mon. functor. Let \mathcal{T}' be the category whose objects are triples (T, U, α) with $T \in \mathcal{T}, U \in \mathcal{U}$, and $\alpha: F(T) \xrightarrow{\sim} U$, and where the morphisms are the obvious ones. Define a product^{functor} in \mathcal{T}' by

$$(T_1, U_1, \alpha_1)(T_2, U_2, \alpha_2) = (T_1 T_2, U_1 U_2, \alpha_1 * \alpha_2)$$

where $\alpha_1 * \alpha_2$ is the isom.

$$F(T_1 T_2) \xrightarrow{\sim} F(T_1) F(T_2) \xrightarrow{\sim} U_1 U_2.$$

~~Utilizing the axioms for a monoidal functors, one verifies easily that this product functor makes \mathcal{T}' into a strict monoidal category.~~ Utilizing the axioms for a monoidal functors, one verifies easily that this product functor makes \mathcal{T}' into a strict monoidal category. Moreover we have



Lemma 0: A monoidal functor $F: \mathcal{S} \rightarrow \mathcal{T}$ is an equivalence of monoidal categories if it is an equivalence of the underlying categories.

Lemma 1: $\lambda: \mathcal{S} \rightarrow \mathcal{S} = \underline{\text{Hom}}^{\mathcal{S}^*}(\mathcal{S}, \mathcal{S})$ is an equivalence of monoidal cats.

with Lemma 2: Let $F: \mathcal{S} \rightarrow \mathcal{U}$ be a mon. functor between strict monoidal categories. Then F may be factored

$$\mathcal{S} \xrightarrow{\mu} \mathcal{S}' \xrightarrow{F'} \mathcal{U}$$

where \mathcal{S}' is a strict mon. cat, F' is a strict mon. functor, and μ is an eq. of mon. cats.

Proof: Let $\lambda: \mathcal{S} \rightarrow \mathcal{S}$ be a mon. functor above, so that we have compatible isos. of mon. functors: $\lambda \circ \lambda = \text{id}$, $\lambda \circ \text{id} = \text{id}$. Let \mathcal{S}' be the cat whose objects are triples (T, U, α) , where $T \in \text{Ob}(\mathcal{S})$, $U \in \text{Ob}(\mathcal{U})$, $\alpha: F\lambda(T) \simeq U$; morphisms in \mathcal{S}' are the obvious ones. Define a product in \mathcal{S}' by the formula

$$(T_1, U_1, \alpha_1)(T_2, U_2, \alpha_2) = (T_1 T_2, U_1 U_2, \alpha_1^* \alpha_2)$$

where $\alpha_1^* \alpha_2$ denotes the isom

$$F\lambda'(T_1 T_2) \simeq F\lambda'(T_1) \cdot F\lambda'(T_2) \xrightarrow{\alpha_1 \cdot \alpha_2} U_1 U_2$$

the former isom being the product isom for the monoidal functor $F\lambda'$. Utilizing the properties

Let $s \rightarrow E_0$ be a cofinal p-pres. functor

Cor: $s^{-1}E_0 \rightarrow s^{-1}E_M \rightarrow s^{-1}E_0$ hegs.

Proof: ~~cofinality~~ cofinality $\Rightarrow s^{-1}E_M$ is an H-space with h-inverse and

$$H_*(s^{-1}E_M) = (\pi_0)^{-1}H_*(E_M) = (\pi_0 s^{-1}E_0)^{-1}H_*(E_M)$$

so ~~it's~~ enough to show (by Whitehead) iso in homology.

Proof of Thm. Have canonical isom

$$E \perp E \simeq E \perp i_* E$$

Enough to use coefficients in the field \mathbb{F}_p, \mathbb{Q} , whence one has a Künneth isom.

~~$$H_*(C_1) \otimes H_*(C_2) \xrightarrow{\sim} H_*(C_1 \times C_2)$$~~

~~It is well-known if s is a cat with~~

$$H_*(X_1) \otimes H_*(X_2) \xrightarrow{\text{iso}} H_*(X_1 \times X_2).$$

This makes $H_*(X)$ into a cogebra with coproduct and counit induced by the diagonal functor

$$X \xrightarrow{\Delta} X \times X$$

and $X \rightarrow \text{pt}$. Also if s is a cat with product $H_*(s)$ is an algebra with product + unit induced by $s \times s \rightarrow s$ and $\text{pt} \rightarrow s$. Recall that we then have a product on the k -module

$$\text{Hom}^{(0)}(H_*(X), H_*(Y))$$

Example: Given $\bar{s} \in \mathcal{S}$, have

$$F(S) = \bar{s}S$$

and $(\bar{s}x)s \simeq \bar{s}(xs)$

In this way we get a monoidal functor

$$\mathcal{S} \longrightarrow \underline{\text{Hom}}^{\delta^0}(\mathcal{S}, \mathcal{S})$$

Assertion: This is an equivalence of monoidal categories.

(need lemma: a monoidal functor is an equiv. of mon. cats \Leftrightarrow it is an equivalence of cats.)

~~All~~ F, φ $\varphi : F(S_1 S_2) \simeq F(\bar{s}_1 s_1, \bar{s}_2 s_2)$.

Then

$$F(0) \times \xrightarrow[\sim]{\varphi} F(0 \times) \simeq F(\times)$$

is an isomorphism of functors of S .

To show it is an isom ~~in~~ in $\underline{\text{Hom}}^{\delta^0}(\mathcal{S}, \mathcal{S})$

$$(F(0) \times)S \longrightarrow F(0 \times)S \longrightarrow F(\times)S$$

$$\downarrow \quad \quad \quad \downarrow$$

$$F(0)(XS) \xrightarrow[S]{} F(0(XS)) \longrightarrow F(XS)$$

and the associativity isomorphism ~~isomorphism~~ gives
an isom

$$h(S_1, S_2) \xrightarrow{\sim} h(S_1) \cdot h(S_2)$$

Also

$$h(O) = \begin{cases} X \xrightarrow{\sim} OX \\ \cancel{O}X^T \xrightarrow{\sim} O(X^T) \end{cases}$$

and unit id. gives $hO \xrightarrow{\sim} \text{id.}$

Now you must verify that these isoms are compatible with the unity data, assoc data of $\text{Hom}^{S^0}(S, S)!!!!$

so you want $O = (M \dashv M)^T$

$$h((S_1, S_2), S_3) \xrightarrow{\sim} h(S_1, S_2)h(S_3)$$

$$h(S, (S_2, S_3))$$

$$h(S_1)h(S_2, S_3) \xrightarrow{\sim} h(S_1)h(S_2)h(S_3)$$

To commutate in $\text{Hom}^{S^0}(S, S)^T$. But two nat. transf. are equal iff equal on objects, so this results from the usual pentagon.

Similarly for

$$h(SO) \xrightarrow{\sim} h(S)h(O)$$

$$(M)^T \leftarrow (hS)^T \leftarrow (M, \dashv M)^T$$

$M \dashv M$ at middle

$$F(0)X \simeq F(0X) \simeq F(X)$$

$$(F(0)X)S \xrightarrow{\varphi} F(0X)S \xrightarrow{\sim} F(X)S$$

$$\begin{array}{c} \text{assoc} \\ \parallel \\ \text{satisfied} \\ \text{by } \varphi \end{array} \quad \begin{array}{c} \varphi \downarrow S \\ \varphi \circ = (\text{nat of } \varphi) \end{array} \quad \begin{array}{c} \varphi \\ \parallel \\ \text{unity} \end{array}$$

$$(F(0)X)S \xrightarrow{\varphi} F(0X)S \xrightarrow{\sim} F(X)S$$

OKAY.

You have a monoidal functor

$$S \rightarrow \underline{\text{Hom}}^S(S, S)$$

To prove this it is necessary ~~to~~

$$S \mapsto \left(\begin{array}{l} X \mapsto SX \\ (SX)T \simeq S(XT) \end{array} \right)$$

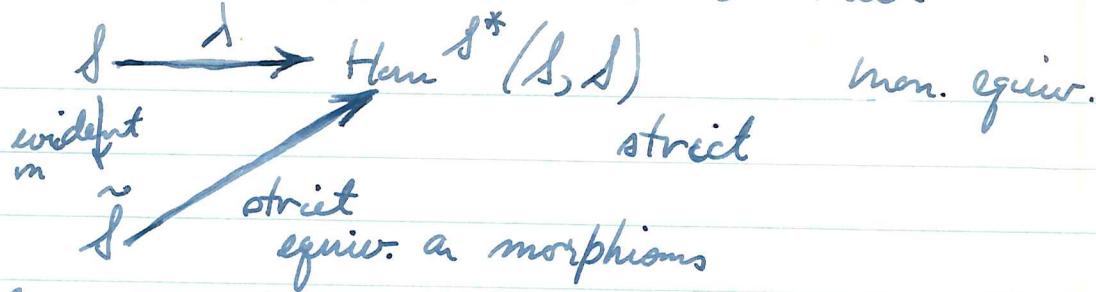
fact this is well defined on objects uses pentagon + unity.

Now compose

$$h(S_1, S_2) = \left\{ \begin{array}{l} X \mapsto \#(S_1 S_2)X \\ ((S_1 S_2)X)T \simeq (S_1 S_2)(XT) \end{array} \right\}$$

$$h_{S_1, S_2} = \left\{ \begin{array}{l} X \mapsto S_1(S_2 X) \\ [S_1(S_2 X)]T \simeq S_1[(S_2 X)T] \simeq S_1(S_2(XT)) \end{array} \right\}$$

so what I know is that I have



and there are base an evident monoidal functors $\delta \rightarrow \tilde{\delta}$
which is a monoidal equivalence

Proof: $\tilde{\delta}$ is an adjoint.

$$\begin{array}{ccc} \delta & \xrightarrow{f} & \mathcal{U} \\ \delta \downarrow & f_1 \nearrow f_2 & \text{---} \\ \tilde{\delta} & \dashrightarrow & \end{array} \quad \mathcal{U} \text{ strict}$$

$$\begin{array}{c} \text{mon. equiv.} \\ \begin{array}{c} f_1(x) \\ f_2(x) \\ \dashv \\ \vdash \\ f(x) \end{array} \end{array}$$

f_1, f_2 strict such that $f_1 \circ \delta = f_2 \circ \delta = f$.

Claim $f_1 = f_2$. Proof: $\text{Ob}(f_i) : \text{Ob}(\tilde{\delta}) \rightarrow \text{Ob}(\mathcal{U})$
two monoid homos. agree on $\text{Ob}(\delta)$ so are equal.

Take an arrow between two objects of $\tilde{\delta}$?

$$\delta \text{ is an equivalence so } \exists! f_1 \simeq f_2 \quad f(x) \simeq f_2(x)$$

$$\begin{array}{ccc} \delta & \xrightarrow{f_0} & \mathcal{U} \\ \delta \downarrow & & \\ \tilde{\delta} & & \end{array} \quad \begin{array}{c} f_1 \circ \delta \simeq f_2 \circ \delta \\ \parallel \quad \parallel \\ f \end{array}$$

$$f_1(x) \simeq f_2(x) \quad ! \text{ natural transf.}$$

Let $\delta' : \tilde{\delta} \rightarrow \delta$ be quasi-inverse.

$$\delta' = \{ (T, U, f_0 \circ \delta' \simeq U) \} \quad \text{strict}$$

and it maps to both $\tilde{\delta}$ and \mathcal{U} strictly

I propose to ~~prove this~~ outline the proof in the following section.

~~PROOF~~ I feel as if the proof should proceed from 2 lemmas.

$$1) \lambda: \mathcal{S} \rightarrow \tilde{\mathcal{S}} = \underline{\text{End}}^{\delta^*}(\mathcal{S}) \quad \text{equivalence}$$

2) If $\mathcal{T} \xrightarrow{\text{strict}} \mathcal{U}$ is a monoidal f. between strict
equiv. \mathcal{G} , \mathcal{H} .

& basic construction

$$\begin{array}{ccc} \mathcal{S} & \xrightarrow{\lambda} & \tilde{\mathcal{S}} \\ \xi \downarrow & \nearrow f & \\ \tilde{\mathcal{S}} & \xrightarrow{\text{st}} & + \text{f. faithful} \end{array} \quad \begin{aligned} \xi(s) &= s \\ \xi(s_1 s_2) &\sim \xi(s_1) \cdot \xi(s_2) \end{aligned}$$

~~Sketch~~

$$\begin{array}{ccc} \mathcal{S} & \xrightarrow{F} & \mathcal{T} \text{ st.} \\ \xi \downarrow & \nearrow F'_1 & \\ \tilde{\mathcal{S}} & \xrightarrow{\quad} & \tilde{\mathcal{F}}'_2 \end{array}$$

$$F'_1 \cdot \xi = F'_2 \cdot \xi \quad \text{as. mon. f.}$$

$$\Rightarrow \exists! \theta: F'_1 \xrightarrow{\sim} F'_2 \quad \text{iso of mon. f.}$$

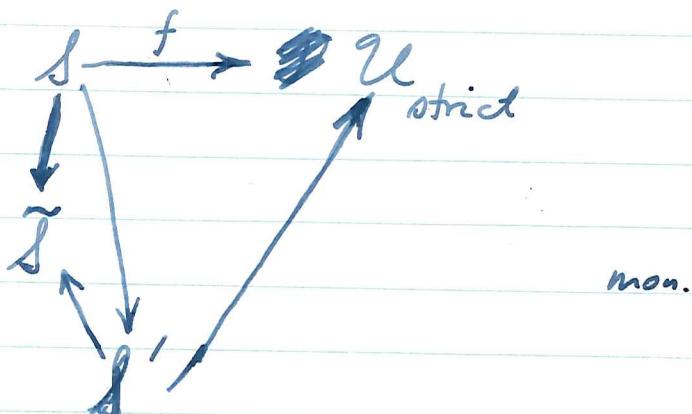
$$\theta \circ F'_1 \xi \simeq F'_2 \xi \quad \text{is id}$$

Consider $\{x \in \text{Ob}(\tilde{\mathcal{S}}) \mid \theta_x: F'_1(x) \simeq F'_2(x) \text{ is an id. map}\}$.

Because θ is an iso. of mon. f. \Rightarrow this set is a submonoid.

$\lambda: \mathcal{S} \rightarrow \tilde{\mathcal{S}}$ equivalence with $\tilde{\mathcal{S}}$ strict

next: let $\tilde{\mathcal{S}} \rightarrow \tilde{\mathcal{T}}$ be strict



make the definitions, then mention standard method.

Proposition: Inclusion

$$\left\{ \begin{array}{l} \text{monoids} \\ \text{in } \text{Cat} \end{array} \right\} \subset \left\{ \begin{array}{l} \text{monoidal} \\ \text{cat + mon. f.} \end{array} \right\}$$

has a left adjoint $S \rightarrow \tilde{S}$

Given a monoidal category S there exists a monoidal f. $S \rightarrow \tilde{S}$ which is a universal mon. functor to

Given a monoidal category S there exists a strict monoidal cat \tilde{S} and a monoidal functor $S \rightarrow \tilde{S}$ which is a universal monoidal functor from S to a strict monoidal cat. in the following sense: Given a mon. f. $S \rightarrow \tilde{T}$ with \tilde{T} strict $\exists!$ str. mon. f. $\tilde{S} \rightarrow \tilde{T}$ such that \tilde{T} comm.

Lemma 2: Let $F: \mathcal{S} \rightarrow \mathcal{U}$ be a mon. functor with \mathcal{U} strict. Then F may be factored

$$\mathcal{S} \xrightarrow{\mu} \mathcal{T} \xrightarrow{F'} \mathcal{U}$$

where ~~is a quasi-inverse of mon. cats.~~ μ is an equiv. of mon. cats., ~~and~~ \mathcal{T} is a strict mon. cat, and F' is a strict mon. functor.

Proof: Let $X: \mathcal{T} \rightarrow \mathcal{S}$ be a quasi-inverse to $\mathcal{S} \rightarrow \mathcal{T}$ above equipped with its canon. mon. structure. ~~Then~~ $FX: \mathcal{T} \rightarrow \mathcal{U}$ is a mon. functor. Let \mathcal{T}' be the cat whose objects are triples (T, u, α) , $T \in \text{Ob } \mathcal{T}$, $u \in \text{Ob } \mathcal{U}$, $\alpha: F(T) \cong u$. with the evident morphisms. Define a product on \mathcal{T}' by

$$(T, u, \alpha)(T', u', \alpha') = (TT', uu', \alpha \circ \alpha')$$

~~Thus it seemed that one has determined δ' .~~

But we have a monoidal functor

$$\delta \rightarrow \delta'$$

$$S \mapsto (\gamma_S, f_0 S, f_0 \delta(\gamma_S) \simeq f_0 S)$$

which is an equivalence of ^{mon.} cats. So we get as above a ~~weak~~ strict

$$\tilde{\delta} \rightarrow \delta'$$

uniqueness Lemma: Let $\delta \xrightarrow{\gamma} \delta'$ be an equiv. of monoidal cats with δ' strict, let $f_1, f_2 : \delta \rightarrow \mathcal{U}$ be strict mon. functors with \mathcal{U} strict and suppose $f_1 \gamma = f_2 \gamma$. Then $f_1 = f_2$, provided $\text{Ob}(\delta)$ generates ^{the mon.} $\text{Ob}(\delta')$.

Proof: Because γ is an equivalence, $\exists!$ natural isom $\theta : f_1 \simeq f_2$ of monoidal functors \exists

$$\theta \cdot r : f_1 \gamma \simeq f_2 \gamma$$

is the identity. ~~Consider~~ Consider $\{x \in \text{Ob}(\delta') \mid \theta_x \text{id}_x\}$. Since

$$\theta_{X_1 X_2} : f_1(X_1 X_2) \simeq f_2(X_1 X_2)$$

$$\quad \quad \quad \parallel \quad \quad \quad \parallel$$

$$f_1(X_1) f_1(X_2) \simeq f_2(X_1) f_2(X_2)$$

$$\theta_{X_1 X_2}$$

follows this set is closed under product, so done.